

Taking Model-Complete Cores

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Finitely Homogeneous Structures

Definition. A relational structure \mathfrak{B} is called

- **homogeneous** if every isomorphism between finite substructures of \mathfrak{B} can be extended to an automorphism of \mathfrak{B} .
- **finitely homogeneous** if it is homogeneous in a finite signature.

Examples.

- $(\mathbb{Q}; <)$
- The Rado graph (aka the Random graph)
- The binary branching countable universal homogeneous C-relation

Facts.

- Finitely homogeneous structures are ω -categorical.
- Henson'1972: there are 2^ω many homogeneous digraphs.
- Cherlin'1998: classification of homogeneous digraphs.

Long-term goal. Classify countable homogeneous structures in a finite relational signature that are **NIP** (i.e., have a **dependent** first-order theory).

Motivation

Many open questions for finitely homogeneous structures \mathfrak{B} :

- 1 **Thomas' conjecture (1991):** \mathfrak{B} has only finitely many first-order reducts, up to interdefinability.
Equivalently: $\text{Sym}(\mathfrak{B})$ has only finitely many closed subgroups that contain $\text{Aut}(\mathfrak{B})$.
- 2 **Macpherson's question 1 (2011):** Does \mathfrak{B} have the **small index property**: every subgroup of $\text{Aut}(\mathfrak{B})$ of index $< 2^\omega$ is open?
- 3 **Macpherson's question 2 (2011):** Does $\text{Aut}(\mathfrak{B})$ have only finitely many closed normal subgroups?
- 4 **Ramsey expansion conjecture:** \mathfrak{B} has an expansion by finitely many relations which is homogeneous and additionally **Ramsey** (B.+Pinsker+Tsankov'2011).
- 5 **CSP dichotomy conjecture:** If $\text{CSP}(\mathfrak{B}) := \{\mathfrak{A} \text{ finite} \mid \mathfrak{A} \rightarrow \mathfrak{B}\}$ is in NP, then $\text{CSP}(\mathfrak{B})$ is in P or NP-complete.
Strengthening of conjecture of B.+Pinsker'2011.
Conjectured for **finite** \mathfrak{B} by Feder+Vardi'1993.
Verified for **finite** \mathfrak{B} by Bulatov and by Zhuk in 2017.

How to classify 1: first-order reducts

Definition. \mathfrak{A} is called **first-order reduct** of \mathfrak{B} if \mathfrak{A} is a reduct of the expansion of \mathfrak{B} by all first-order definable relations.

Example. $(\mathbb{Q}; \{(x, y, z) \mid x < y < z \vee z < y < x\})$ is first-order reduct of $(\mathbb{Q}; <)$.

Definition. Two structures are **first-order interdefinable** if they are first-order reducts of each other.

Example. $(\mathbb{Q}; \{(x, y, z) \mid x < y \vee x < z\})$ is interdefinable with $(\mathbb{Q}; <)$.

Fact. If $\mathfrak{A}, \mathfrak{B}$ are ω -categorical, then:

\mathfrak{A} and \mathfrak{B} are interdefinable **if and only if** $\text{Aut}(\mathfrak{A}) = \text{Aut}(\mathfrak{B})$

(follows from theorem of Ryll-Nardzewski, Engeler, Svenonius).

Warning. Finite homogeneity **not** preserved by first-order reducts!

How to classify 2: first-order interpretations

Definition. A (d -dimensional, first-order) interpretation of \mathfrak{A} in \mathfrak{B} is a partial surjective map $I: A^d \twoheadrightarrow B$ such that the pre-images of definable sets in \mathfrak{B} are definable in \mathfrak{A} .

Example. The following structure is first-order definable in $(\mathbb{Q}; <)$:

$$(\mathbb{Q}^2; <_1, <_2) \text{ with } <_i := \{((a_1, a_2), (b_1, b_2)) \mid a_i < b_i\}$$

Definition. \mathfrak{A} and \mathfrak{B} are bi-interpretable if there are interpretations $I: \mathfrak{A}^{d_1} \twoheadrightarrow \mathfrak{B}$ and $J: \mathfrak{B}^{d_2} \twoheadrightarrow \mathfrak{A}$ such that $I \circ J$ is definable in \mathfrak{B} and $J \circ I$ is definable in \mathfrak{A} .

Example. The structure $(\mathbb{Q}; =)$ is bi-interpretable with

$$(\mathbb{Q}^2; \{((a_1, a_2), (b_1, b_2)) \mid a_2 = b_1\})$$

$$\text{and with } \left(\binom{\mathbb{Q}}{2}; \{(A, B) : |A \cap B| = 1\}\right)$$

Fact (Ahlbrand-Ziegler'86, Coquand): ω -categorical structures \mathfrak{A} and \mathfrak{B} are bi-interpretable if and only if $\text{Aut}(\mathfrak{A})$ and $\text{Aut}(\mathfrak{B})$ are topologically isomorphic.

The Stable Case

Stability: Shelah 1969.

Definition. An ω -categorical structure \mathfrak{M} is **stable** if and only if there is no formula $\phi(\bar{x}, \bar{y})$ and parameters $a_0, a_1, \dots \in M^{|\bar{x}|}$ and $b_0, b_1, \dots \in M^{|\bar{y}|}$ such that $\mathfrak{M} \models \phi(a_i, b_j)$ if and only if $i < j$.

Examples. $(\mathbb{N}; =)$.

The countable vector space over a finite field (not finitely homogeneous)

Non-examples.

- $(\mathbb{Q}; <)$
- the Rado graph

Fact. Stability (and ω -categoricity) preserved by first-order interpretations.

Examples. $(\mathbb{Q}^2; \{((a_1, a_2), (b_1, b_2)) \mid a_2 = b_1\})$

and $(\binom{\mathbb{Q}}{2}; \{(A, B) : |A \cap B| = 1\})$ are stable and ω -categorical.

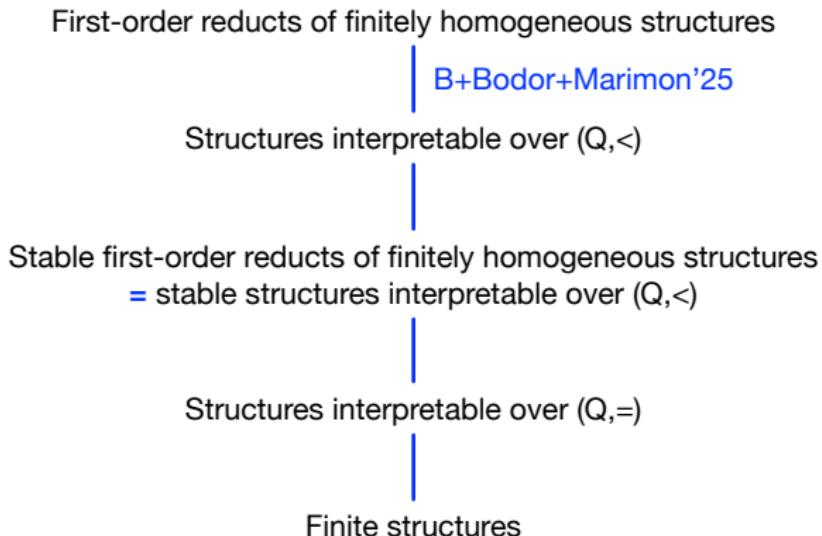
Lachlan's Theorem

Theorem (Lachlan'1986) The following are equivalent.

- \mathcal{B} is stable and a first-order reduct of a finitely homogeneous structure;
- \mathcal{B} is stable and interpretable over $(\mathbb{Q}; <)$.

'Lachlan's class'.

Picture:



Contributions, Part 1

B.+Bodor+Marimon'25:

- 1 Conjecture: Lachlan's class equals closure of class of structures interpretable over $(\mathbb{Q}; =)$ under taking model-complete cores.
- 2 Taking model-complete cores preserves most model-theoretic properties (e.g., stability).
- 3 Show that class of structures interpretable over $(\mathbb{Q}; =)$ is not closed under taking model-complete cores.
- 4 Consequence (was known before, but not explicitly in the literature): Lachlan's class is distinct from class of structures interpretable over $(\mathbb{Q}; =)$.

How to classify 3: model-complete cores

An ω -categorical structure \mathfrak{B} is called a

- **core** if every endomorphism $\mathfrak{B} \rightarrow_{\text{hom}} \mathfrak{B}$ is an embedding,
- **model-complete** if every embedding $\mathfrak{B} \hookrightarrow \mathfrak{B}$ preserves all first-order formulas.

Fact (B.+Junker'09). \mathfrak{B} is a model-complete core

if and only if $\overline{\text{Aut}(\mathfrak{B})} = \text{End}(\mathfrak{B})$.

\mathfrak{A} and \mathfrak{B} are **homomorphically equivalent** if $\mathfrak{A} \rightarrow_{\text{hom}} \mathfrak{B}$ and $\mathfrak{B} \rightarrow_{\text{hom}} \mathfrak{A}$.

Theorem (B.'2005, B.+Hils+Martin'2010). Every ω -categorical \mathfrak{B} is homomorphically equivalent to a model-complete core \mathfrak{C} , which is unique up to isomorphism, and ω -categorical.

Theorem (B.+Bodor+Marimon'26). If \mathfrak{B} is stable, then so is \mathfrak{B} .

Similarly, taking model-complete cores preserves NIP, NSOP, simplicity, superstability, monadic stability, monadic NIP, strong minimality, etc.

Examples

- The model-complete core of $(\mathbb{Q}; \leq)$ is a single point.
- The model-complete core of the Random Graph is K_ω .
- The model-complete core of $(\mathbb{Q} \cap [0, 1]; <)$ is $(\mathbb{Q}; <)$.
- The model-complete core of $(\mathbb{Q}^2; <_1, <_2)$ is the countable universal homogeneous permutation structure (aka [generic permutation structure](#)).
- Similarly: Cherlin's tournament [\$S\(2\)\$](#) is a model-complete core of a structure interpretable over $(\mathbb{Q}; <)$.

The Challenge

Goal: find a structure with is interpretable over $(\mathbb{Q}; =)$, but whose model-complete core is not interpretable over $(\mathbb{Q}; =)$.

Typical approach: To prove that a structure \mathfrak{A} is not interpretable over \mathfrak{B} :

- Identify a property P that is preserved by first-order interpretations,
- Show that \mathfrak{B} has property P ,
- Show that \mathfrak{A} does not have property P .

Challenge: almost all model-theoretic properties are preserved by taking model-complete cores!

The Counterexample

$$X := \{(a, b, m) : a, b \in \mathbb{Q}, m \in \mathbb{Z}_4, a \neq b\}.$$

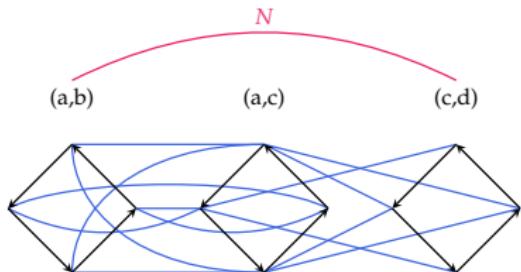
\mathfrak{X} : structure with domain X and relations

- $R := \{((a, b, m), (a, b, m+1)) \mid a, b \in \mathbb{Q}, m \in \mathbb{Z}_4\}$,
- $E \subseteq X^2$, the binary relation given by

$$\begin{aligned} & \{((a, b, 2m), (a, c, 2n)) : n, m \in \mathbb{Z}_4, b \neq c\} \\ & \cup \{((a, b, 2m), (c, a, 2n+1)) : n, m \in \mathbb{Z}_4, b \neq c\} \\ & \cup \{((b, a, 2m+1), (a, c, 2n)) : n, m \in \mathbb{Z}_4, b \neq c\} \\ & \cup \{((b, a, 2m+1), (c, a, 2n+1)) : n, m \in \mathbb{Z}_4, b \neq c\}. \end{aligned}$$

- the binary relation

$$N := \{((a, b, m), (c, d, n)) : n, m \in \mathbb{Z}_4, \{a, b\} \cap \{c, d\} = \emptyset\}.$$



Verifying the counterexample

- 1 \mathfrak{X} is interpretable over $(\mathbb{Q}; =)$.
- 2 the substructure \mathfrak{Y} induced by $Y := \{(a, b, m) : a, b \in \mathbb{Q}, m \in \mathbb{Z}_4, a < b\}$ is homomorphically equivalent to \mathfrak{X} .
- 3 \mathfrak{Y} is interpretable over $(\mathbb{Q}; <)$.
- 4 the structure \mathfrak{Y} is a model-complete core.
- 5 \mathfrak{Y} is **not** interpretable over $(\mathbb{Q}; =)$:
 - every structure \mathfrak{A} with a first-order interpretation in $(\mathbb{Q}; =)$ is finite or $\text{Aut}(\mathfrak{A})$ has a subgroup isomorphic to $\text{Sym}(\mathbb{Q})$.
 - The subgroup of involutions of $\text{Aut}(\mathfrak{Y})$ is abelian.

Beyond Stability

$(\mathbb{Q}; <)$ is not stable, but **NIP**.

Definition. An ω -categorical structure \mathfrak{B} has the **Independence Property (IP)** if there exists a formula $\phi(\bar{x}, \bar{y})$ and an infinite $S \subseteq B^{|x|}$ and $b_S \in B^{|y|}$ for every finite $I \subseteq S$ such that

$$\mathfrak{B} \models \phi(a, b_i) \text{ if and only if } a \in I.$$

Fact (Shelah): every ω -categorical unstable structure \mathfrak{B} has IP or the **SOP**: there exists a formula $\phi(\bar{x}, \bar{y})$ and $a_0, a_1, \dots \in B^{|y|}$ such that

$$\phi(\bar{x}, b_0)^\mathfrak{B} \subsetneq \phi(\bar{x}, b_1)^\mathfrak{B} \subsetneq \dots$$

Examples:

- $(\mathbb{Q}; <)$, $S(2)$, and the generic permutation have NIP and SOP.
- The Rado graph and the Henson digraphs have IP and NSOP.

Structures Interpretable over $(\mathbb{Q}; <)$

Definition. binary finitely homogeneous structure \mathfrak{B} :
all relations of \mathfrak{B} have arity at most 2.

Example. The countable universal homogeneous binary branching C -relation

- is finitely homogeneous and NIP,
- but is not first-order reduct of binary finitely homogeneous structure.

Fact (B+Bodor+Marimon'25). Every structure with a first-order interpretation in $(\mathbb{Q}; <)$ is

- the reduct of a binary finitely homogeneous (Ramsey) structure,
- NIP.

Note. All the mentioned conjectures are already open for NIP structures that are reducts of binary finitely homogeneous structures.

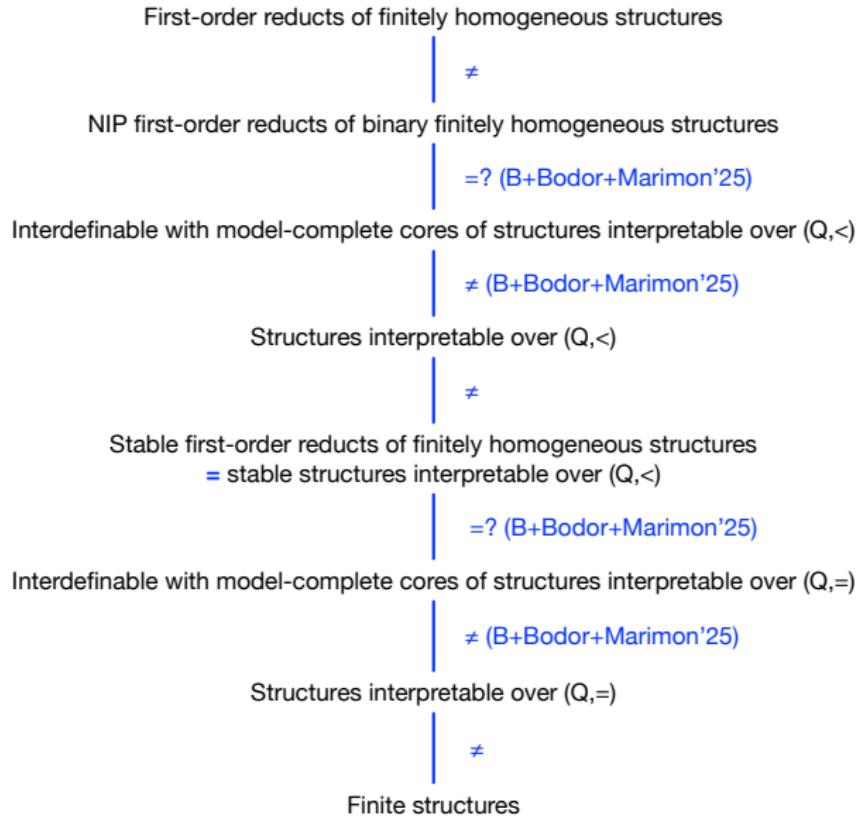
Binary finitely homogeneous NIP structures

Theorem (B+Bodor+Marimon'25). The class \mathcal{I} of structures with a first-order interpretation in $(\mathbb{Q}; <)$ is **not** closed under taking model-complete cores:

- the countable universal homogeneous permutation is not in \mathcal{I}
- $S(2)$ is not in \mathcal{I} .

Conjecture (B+Bodor+Marimon'25). Every NIP structure which is a first-order reduct of a binary finitely homogeneous structure is interdefinable with the model-complete core of a structure interpretable over $(\mathbb{Q}; <)$.

Fuller Picture



Concluding Remarks

- Two types of results in B+Bodor+Marimon'25:
 - Preservation of model-theoretic properties under taking model-complete cores.
 - Techniques to show that certain structures do *not* have a first-order interpretation in another structure.
 - We believe that model-complete cores are a useful concept in classifications of finitely homogeneous structures.
- Most definitions and results about model-complete cores do not rely on ω -categoricity.
- Related work:
 - **positive model theory** (see, e.g., Ben-Yaacov'2003, Dmitrieva+Gallinari+Kamsma'23)
 - **trace definability** by Walsberg'2025.

Fact. The model-complete core of an ω -categorial structure \mathfrak{A} has a trace definition in \mathfrak{A} .