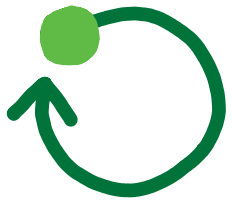


The status of a smooth digraph:
the first hardness criterion for infinite
directed graph-coloring problems



Johns Bruner
TU Wien



joint work with M. Kozik, T. Nguyen, M. Pinski



ÖAW

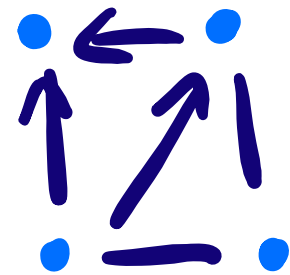
ÖSTERREICHISCHE
AKADEMIE DER
WISSENSCHAFTEN

Boolean Algebra & Logic Seminar Spring '25

ON THE ROAD TO THE CSP-DICHOTOMY

$A = (A; (R_i)_{i \in I})$ REL.

STRUCTURE WITH SIGNATURE \mathcal{J}



$CSP(A)$:

- INPUT: X FINITE \mathcal{J} -STRUCTURE
- QUESTION: DOES THERE EXIST A HOMOMORPHISM $X \rightarrow A$

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EXAMPLES:

(i) $\text{CSP}(\triangle) = 3\text{-COLOURING}$

CSP(A):

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EXAMPLES:

(i) $\text{CSP}(\triangle) = 3\text{-COLOURING}$

(ii) $\text{CSP}(\mathbb{Q}; <) = \text{DIGRAPH ACYCLICITY}$

CSP(A):

- INPUT: X FINITE \mathcal{J} -STRUCTURE
- QUESTION: DOES THERE EXIST A HOMOMORPHISM $X \rightarrow A$

EXAMPLES:

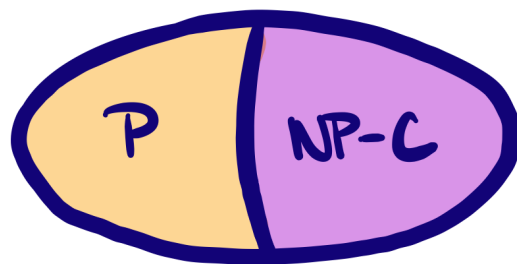
(i) $\text{CSP}(\{\triangle\}) = 3\text{-COLOURING}$

(ii) $\text{CSP}(\{0, <\}) = \text{DIGRAPH ACYCLICITY}$

(iii) $\text{CSP}(\{Z, \{0\}, \{1\}, +, \cdot\}) = \text{DIOPHANTINE EQU.}$

DICHOTOMY-THM (BULATOV, ZHUK '17)

A FINITE. THEN $CSP(A)$ IS
TRACTABLE OR NP-COMPLETE.



'78

'33

'17 

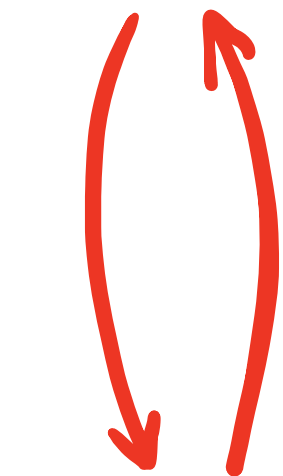
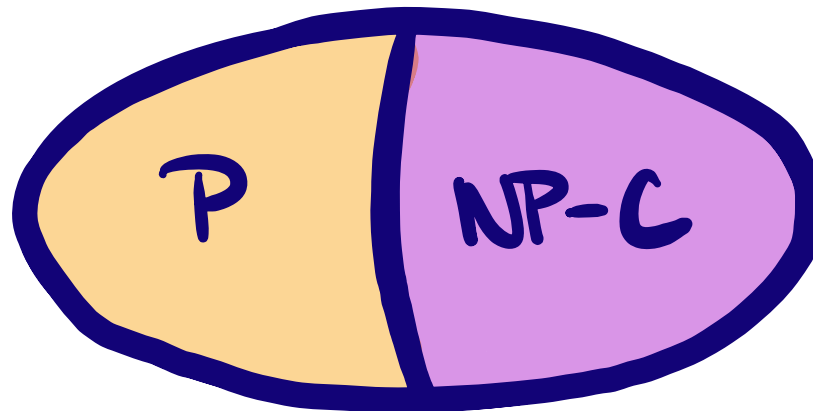
Schaefer:
2-d domain

Feder +
Vardi
conjecture

Bulatov +
Zhuk
proof

STRUCTURAL DICHOTOMIES

CSP(A) ...

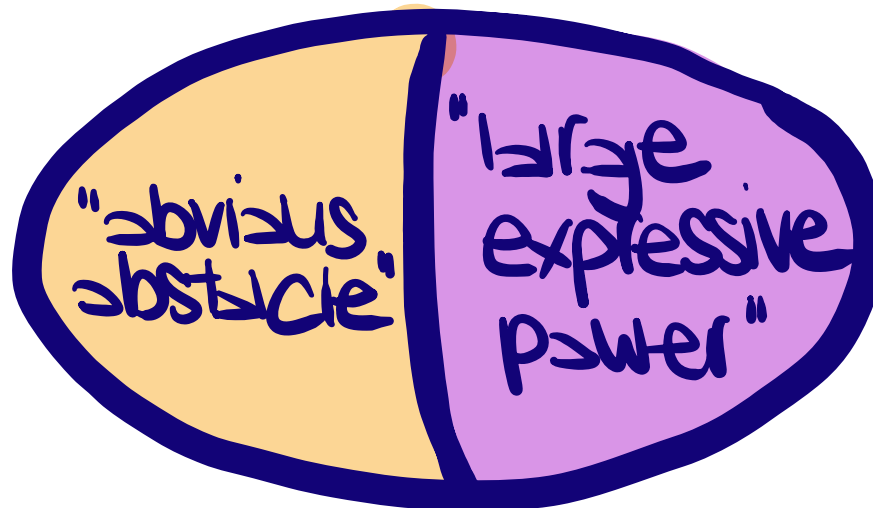
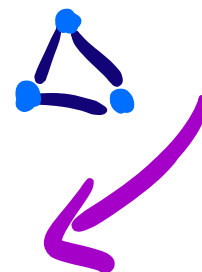


A ...

POLYMORPHISMS,
LOOPS



CONSTRUCT
STH HARD





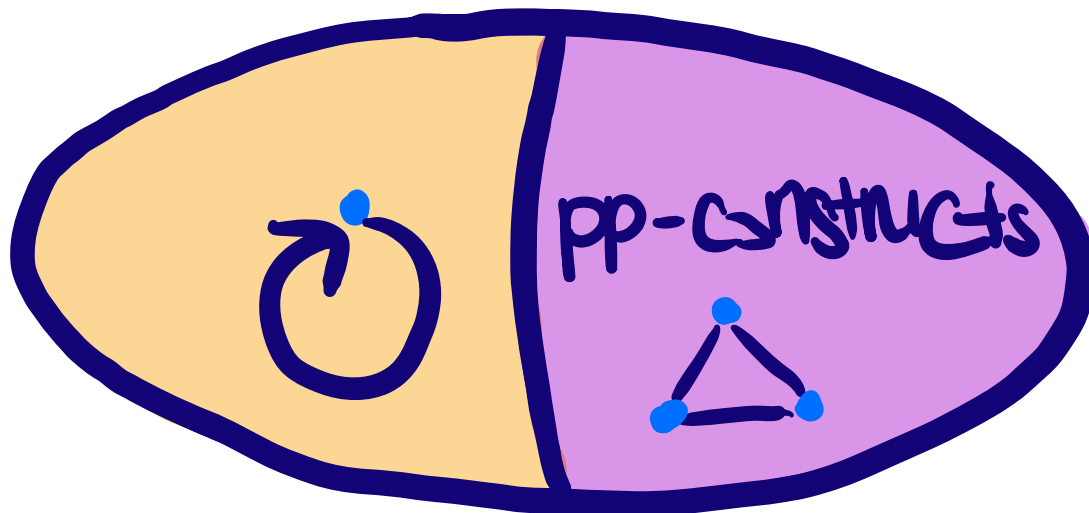
LOOP LEMMATA

Thm (HELL+NEŠETŘIL '90)

⊆ FINITE **UNDIRECTED**
NON-BIPARTITE.

Thm (BARTO+KOZIK+NIEN '03)

⊆ FINITE **DIRECTED** SMOOTH
ALGEBRAIC LENGTH 1.

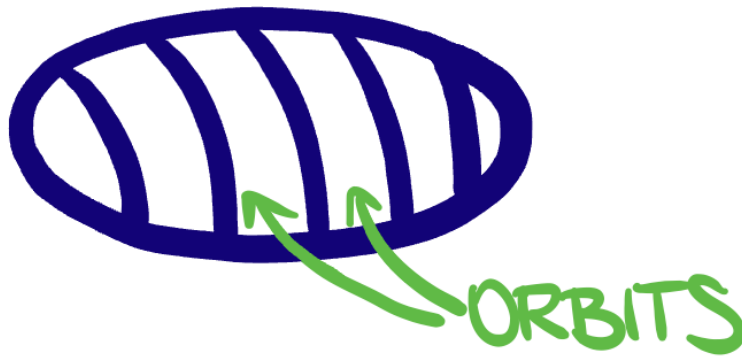


INFINITY

$$\Omega \leq \text{Aut}(\mathbb{E})$$

- \mathbb{E}/Ω : FACTOR GRAPH MOD 1- Ω -ORBIT EQUIV.
I.E. $O \rightarrow P$ IFF $\exists x \in O, y \in P: x \rightarrow y$ IN \mathbb{E}
- Ω IS OLIGOMORPHIC IFF $\forall k: G^k/\Omega$ FINITE

$$\Omega \cong G^k$$



EXAMPLE:
($\mathbb{Q}_i <$)

"CLOSE TO FINITE: 1-ORBITS
TAKE ROLE OF ELEMENTS"

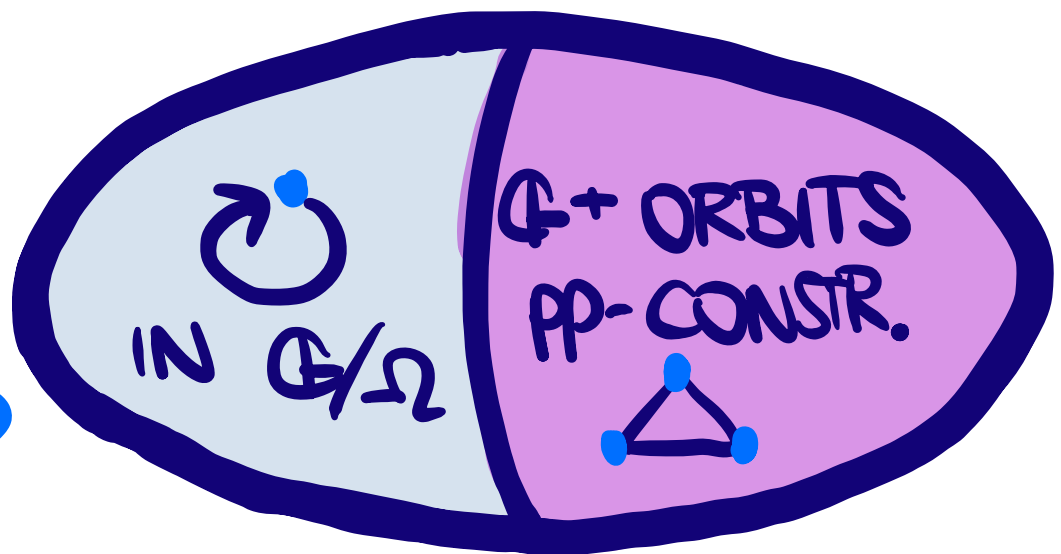
LIFTING FINITE TO INFINITE

HELL-NEŠETŘIL LIFTED:

Thm (BARTO + BODOR + KOZIK + MOTTET + PINSKER '23)

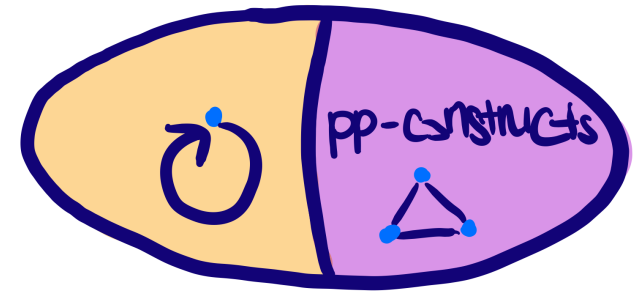
\mathbb{Q} SMOOTH, $\Omega \subseteq \text{Aut}(\mathbb{Q})$ OLIGOMORPHIC,
 \mathbb{Q}/Ω SYMMETRIC, NON-BIPARTITE.

"PSEUDO-
LOOP"
WRT. Ω



OUR DREAM

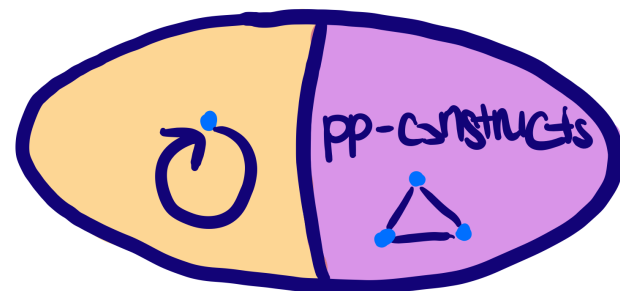
Thm (BARTO+KOZIKINEN '03)
① FINITE DIRECTED SMOOTH
ALGEBRAIC LENGTH 1.



CAN WE LIFT THIS TO THE INFINITE?

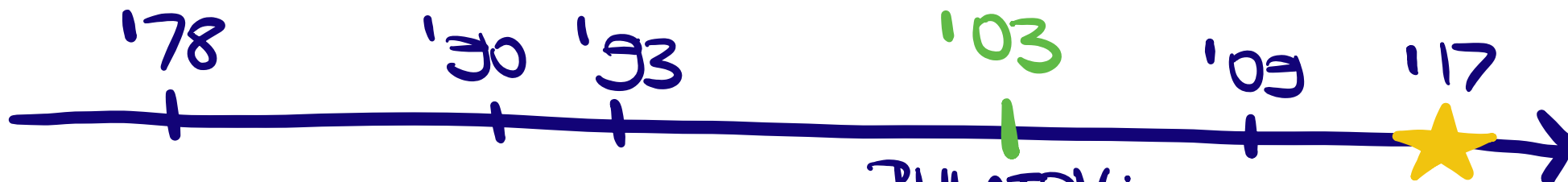
OUR DREAM

Thm (BARTO+KOZIKINEN '03)
⊆ FINITE DIRECTED SMOOTH
ALGEBRAIC LENGTH 1.



CAN WE LIFT THIS TO THE INFINITE?

CONSERVATIVITY



BULATOV:

CONSERVATIVE TEMPATES, I.E.
ALL SUBSETS ARE UNARY RELS.

CONSERVATIVE GRAPH-COLOURING (= LIST HOMOMORPHISM)

LHOM (\mathbb{G}):

INPUT: FINITE DIGRAPH H , $L: H \rightarrow \mathcal{P}(\mathbb{G})$

QUESTION: DOES THERE EXIST A HOMOMORPHISM
 $h: H \rightarrow \mathbb{G}$ WITH $h(x) \in L(x) \quad \forall x \in H$?

$\Omega \leq \text{AUT}(\mathbb{G})$ **OLIGOMORPHIC** \Rightarrow 1-ORBITS
OF Ω TAKE ROLE OF ELEMENTS!

K-CONSERVATIVE EXPANSION OF \mathbb{G} WRT. Ω :

\mathbb{G} EXPANDED BY ALL UNARY RELS
 $O_1 \cup \dots \cup O_k$ WHERE O_i IS 1-ORBIT OF Ω

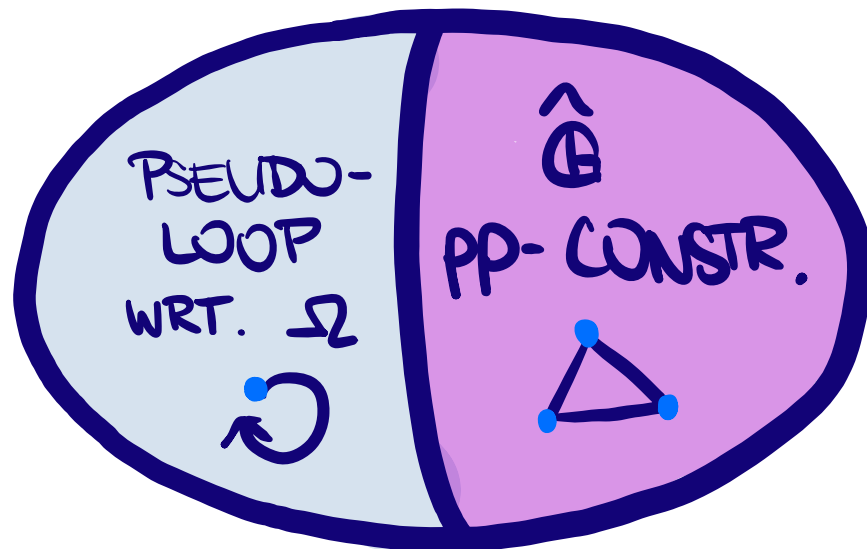
OUR RESULT: THE FIRST STRUCTURAL DICHOTOMY FOR INFINITE DIGRAPHS

Thm (KOZIK+NAGY+PINSKER+B. '25)

\mathbb{G} SMOOTH DIGRAPH OF ALGEBRAIC LENGTH 1
 $\Omega \subseteq \text{Aut}(\mathbb{G})$ OLI EOMORPHIC.

$\hat{\mathbb{G}}$:= 2-CONSERVATIVE EXPANSION OF \mathbb{G} WRT. Ω

STRUCTURAL
DICHOTOMY:

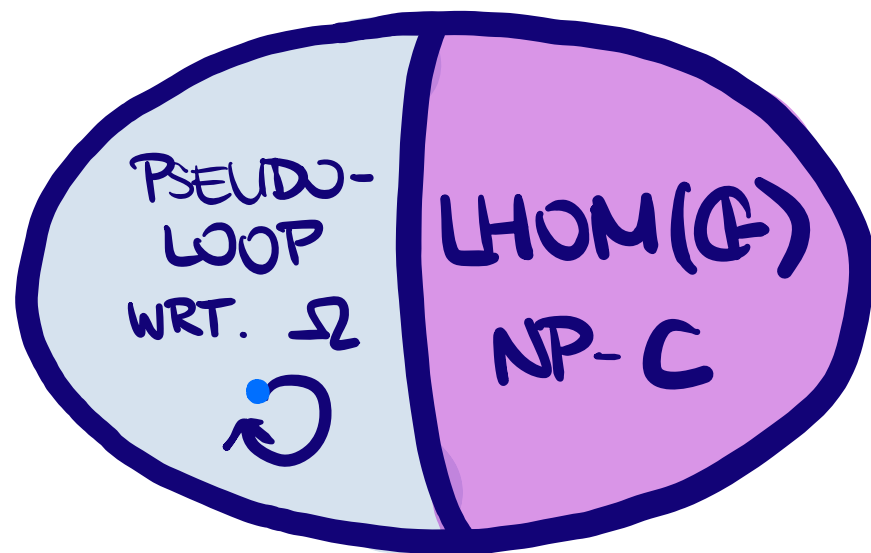


OUR RESULT: THE FIRST HARDNESS
CRITERION FOR INFINITE DIGRAPHS

\mathbb{G} (KOZIK + NAGY + PINSKER + B. '25)

\mathbb{G} SMOOTH DIGRAPH OF ALGEBRAIC LENGTH 1
 $\Omega \subseteq \text{Aut}(\mathbb{G})$ OLIGOMORPHIC.

HARDNESS
CRITERION:

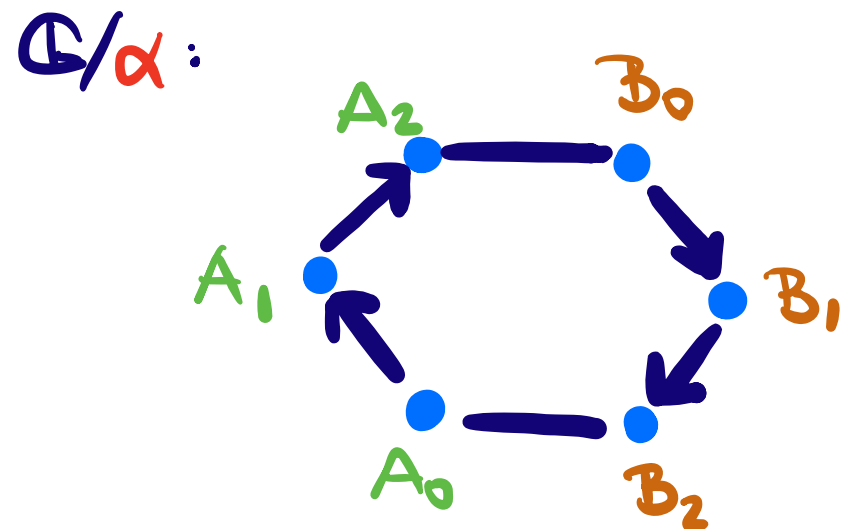
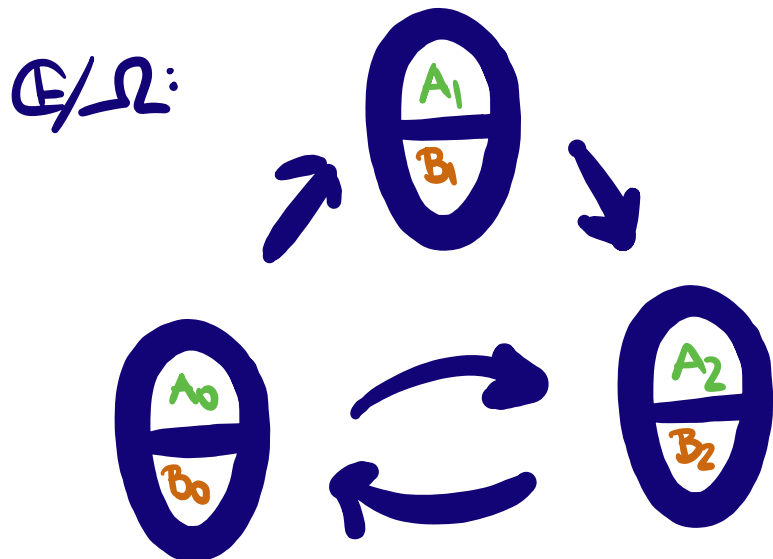


FINITISING \mathbb{G} WRT. Ω

$$\mathbb{G} = (G_i \rightarrow), \quad \Omega \leq \text{Aut}(\mathbb{G})$$

"APPLY ESTABLISHED TECHNIQUES TO FACTOR $\text{mod } \alpha$ "

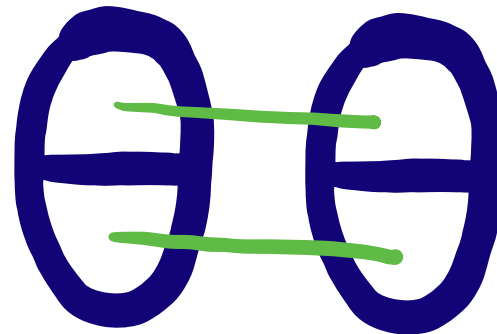
- α IS REFINEMENT OF 1- Ω -ORBIT-EQUIVALENCE
- \mathbb{G}/α FINITE
- α IS Ω -INVARIANT
- \rightarrow IS BIJECTION BETWEEN α -BLOCKS



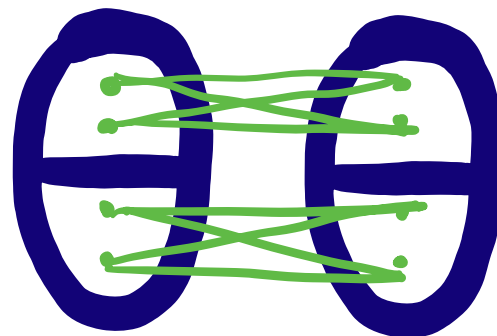
HOW TO LIFT RESULTS FROM \mathbb{C}/α TO \mathbb{C} ?

BLOW-UP RELATIONS IN FACTOR:

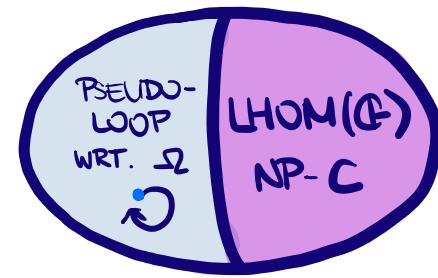
RELATION
ON \mathbb{C}/α :



α -BLOW-UP:



THE MASTER PROOF



SPS : G DOES NOT HAVE PSEUDOLOOP WRT. Ω

SOURCE OF HARDNESS:

$OR(\sigma, \sigma) \equiv \sigma(x, y) \vee \sigma(z, w)$ FOR σ EQUIV.
ON A FINITE SET.

IDEA:



- DEFINE α -BLOW-UP OF $OR(\sigma, \sigma)$
ON SUBSET OF G/α
- OTHERWISE, REDUCE TO SUITABLE
SUBGRAPH H

THE MASTER PROOF

INDUCTION ON NUMBER OF α -BLOCKS

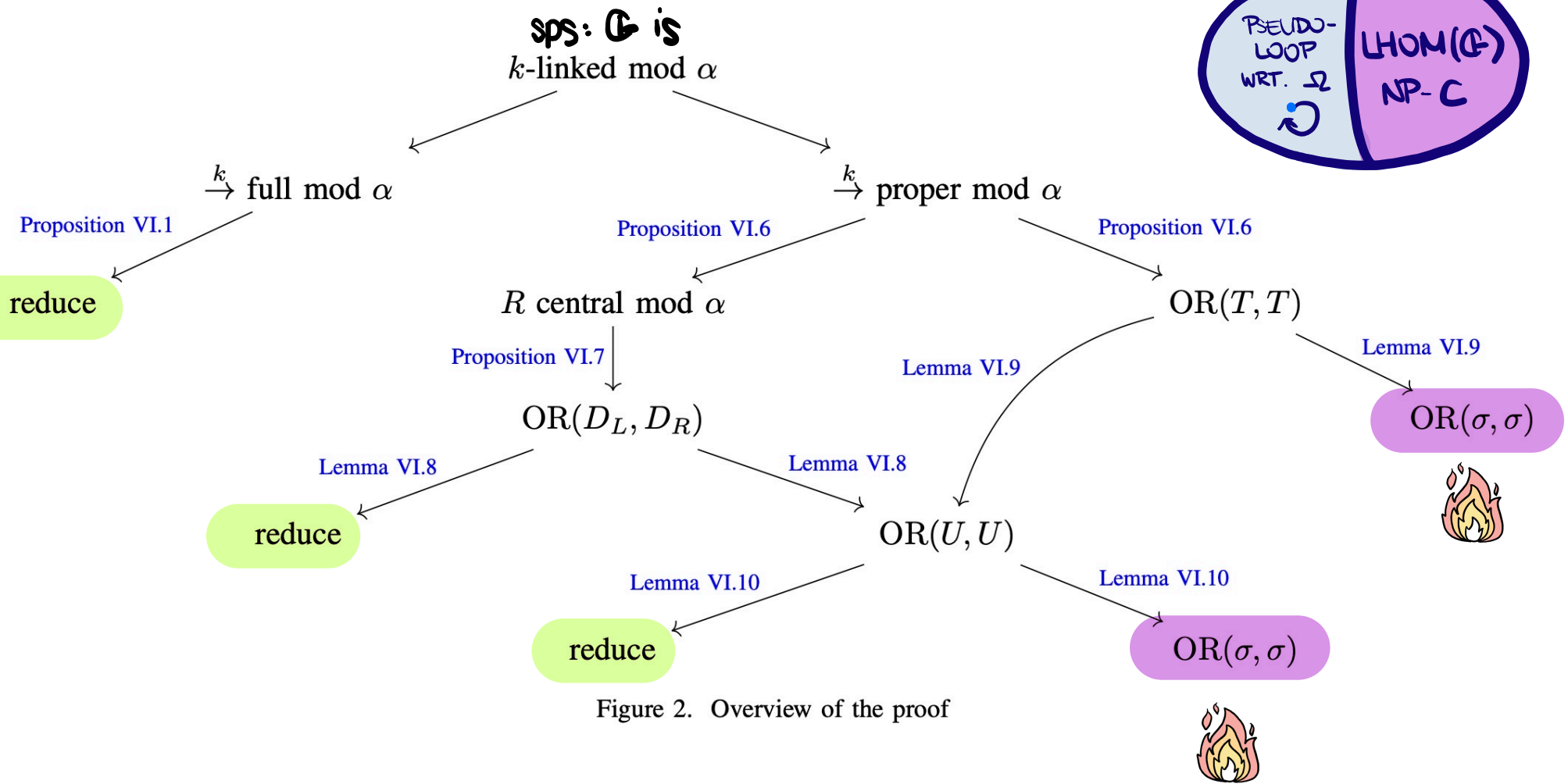
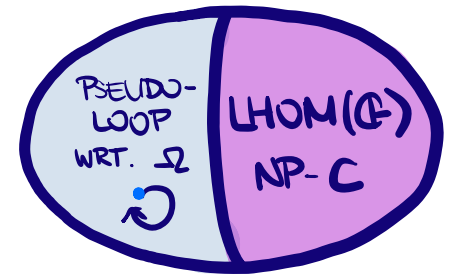
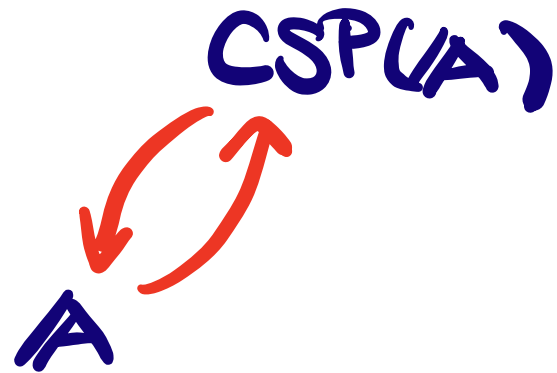
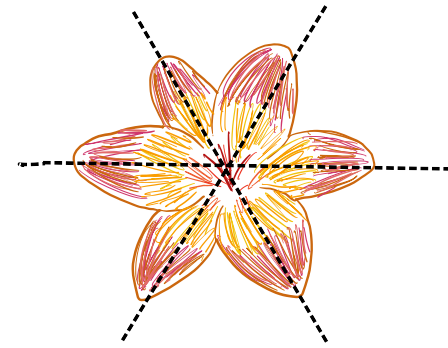
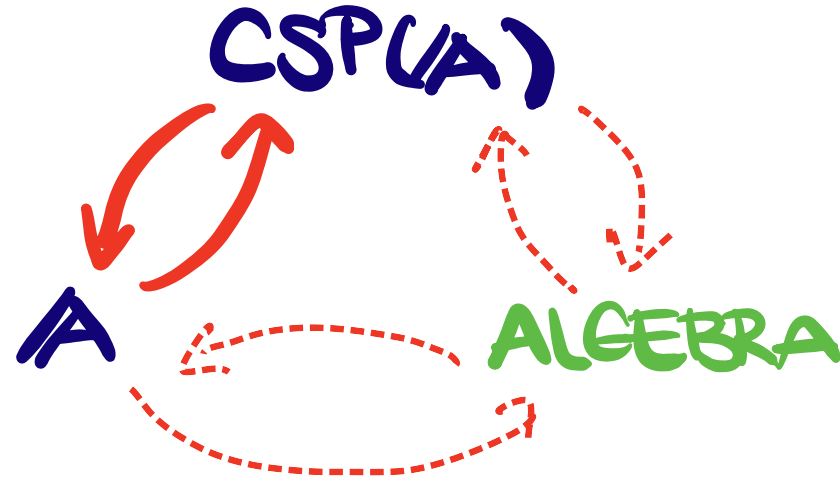
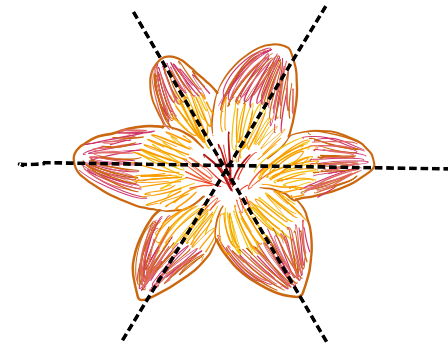


Figure 2. Overview of the proof

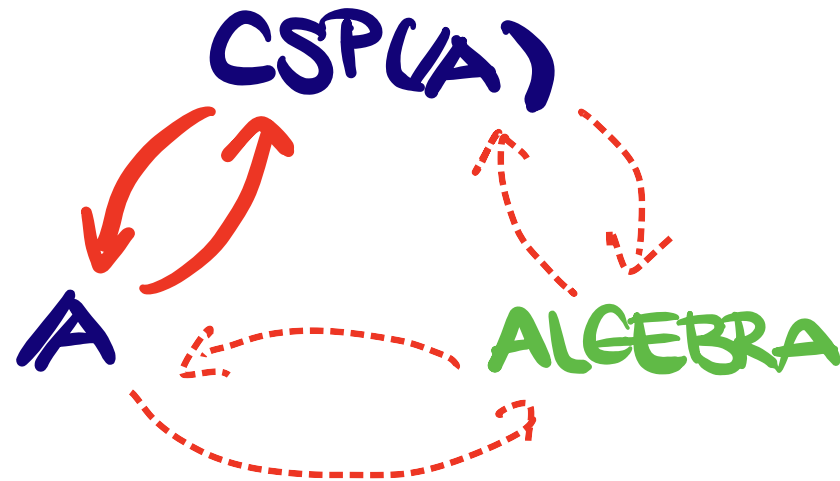
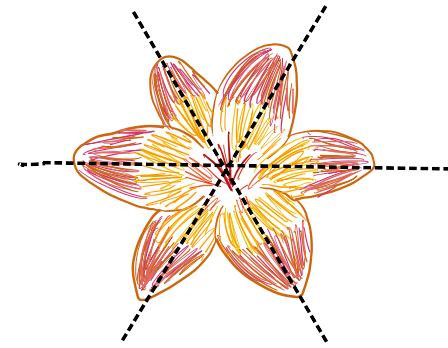
ALGEBRAIC CONSEQUENCES



ALGEBRAIC CONSEQUENCES



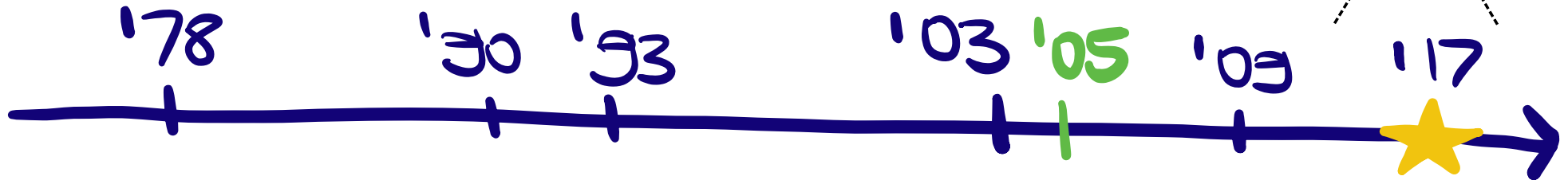
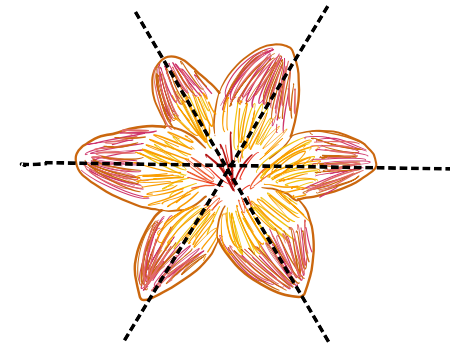
ALGEBRAIC CONSEQUENCES



$$P_{\perp}(A) = \bigcup_{n \geq 1} \text{Hom}(A^n, A)$$

WHAT IDENTITIES ARE SATISFIED BY OPERATIONS FROM $P_{\perp}(A)$?

ALGEBRAIC CONSEQUENCES

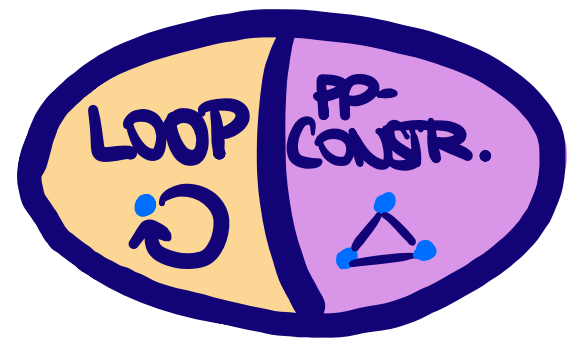


TM (BULATOV + HEAVONS + KROKHIN '05)

COMPLEXITY OF $CSP(A)$ IS COMPLETELY
DETERMINED BY THE EQUATIONAL
VARIETY GENERATED BY $PI(A)$

⇒ ALGEBRAIC APPROACH TO CSP

LOOP LEMMATA

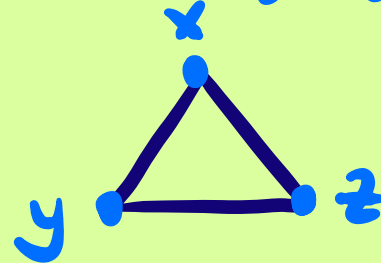


IDENTITIES SATISFIED BY OPERATIONS IN $\text{P}_1(A)$ CAN BE DERIVED FROM LOOPS IN DIGRAPHS:

Thm: A FINITE, ORBITS PP-DEFINABLE. $\text{P}_1(A) \equiv$

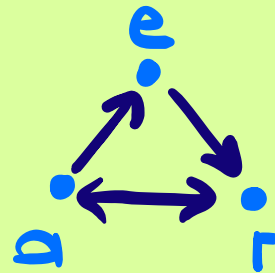
• A DOES NOT PP-CONSTRUCT 

• $\text{P}_1(A) \equiv s(x, y, x, z, y, z) \approx s(y, x, z, x, z, y)$



SIGGERS
(6-ARY)

• $\text{P}_1(A) \equiv s(\triangleright, r, e, \triangleright) \approx s(r, \triangleright, r, e)$



SIGGERS
(4-ARY)

• $\text{P}_1(A) \equiv w(y, x \dots x) \approx w(x, y, x \dots x) \dots w(x \dots x, y)$ WN4

PSEUDOLOOP - LEMMATA

PSEUDOLOOPS GIVE PSEUDO-IDENTITIES!

6-ARY PSEUDO-SIGGERS (BARTO+PINSKER '16)

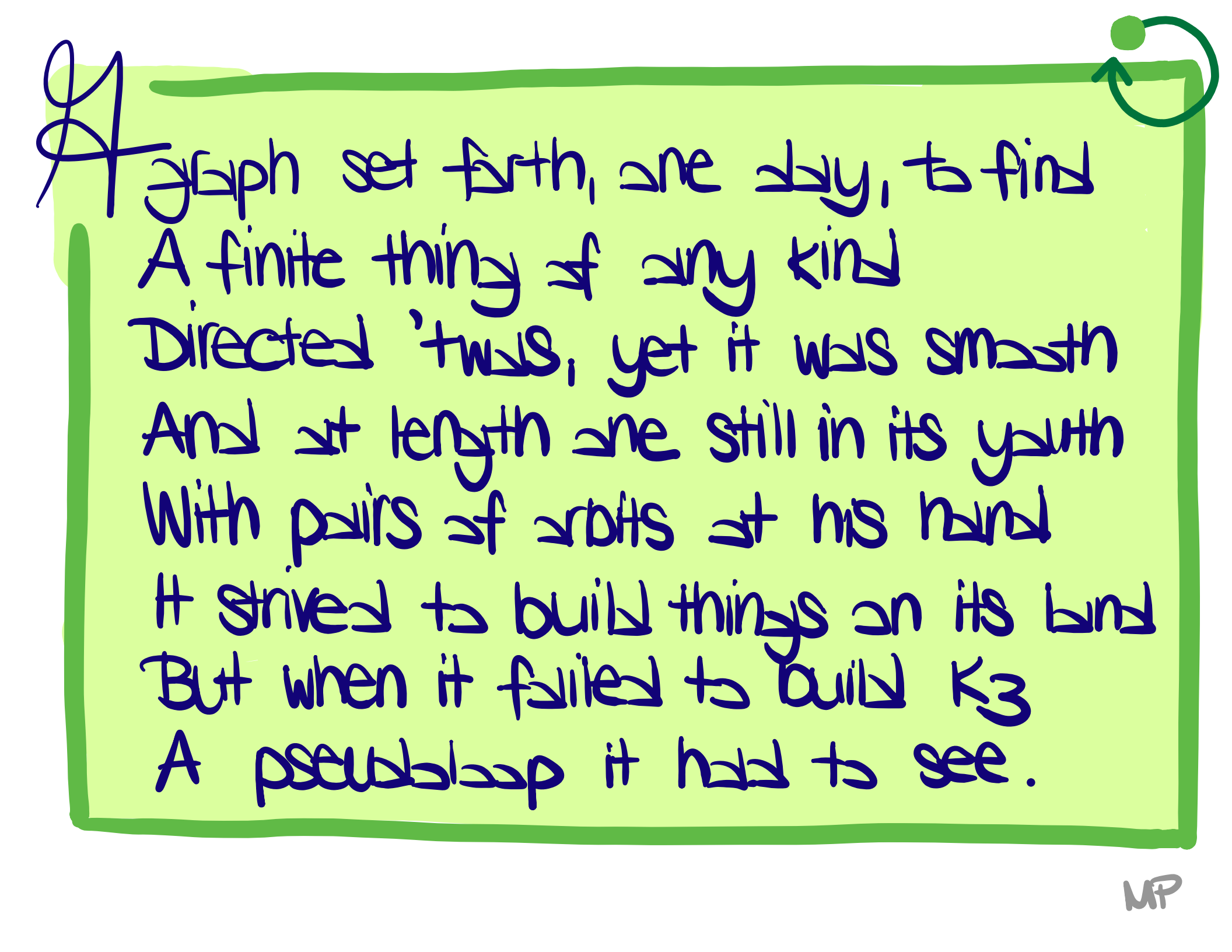
$$u \circ s(x, y, x, z, y, z) \approx v \circ s(y, x, z, x, z, y)$$

Cor (KOZIK+NAGY+PINSKER+B. '25)

Ω OLIGOMORPHIC, \hat{A} EXPANSION OF A BY ALL PUQ FOR P, Q k -ORBITS. EITHER

- \hat{A} PP-CONSTRUCTS  OR

- $\mathcal{B}_1(\hat{A}) = u \circ s(\triangleright, r, e, \triangleright) \approx v \circ s(r, \triangleright, r, e)$



A graph set forth, one day, to find
A finite thing of any kind
Directed 'twas, yet it was smooth
And at length one still in its youth
With pairs of roots at his hand
It strived to build things on its band
But when it failed to build K_3
A pseudobloop it had to see.