

Minimal operations over permutation groups

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May 31st, 2024



TECHNISCHE
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POCOCOP ERC Synergy Grant No. 101071674

Outline

- ① Minimal (and almost minimal) operations
- ② Results
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Operation clones and the pointwise convergence topology

Definition 1 (Operation clone)

Let B be a (possibly infinite) set, $\mathcal{O}^{(n)} = B^{B^n}$ and $\mathcal{O} := \bigcup_{n \in \mathbb{N}} \mathcal{O}^{(n)}$. We call $\mathcal{C} \subseteq \mathcal{O}$ an **operation clone** over B if

- \mathcal{C} contains all projections;
- \mathcal{C} is closed under composition.

We equip \mathcal{O} with the **pointwise convergence topology**: For $\mathcal{S} \subseteq \mathcal{O}$, $f \in \overline{\mathcal{S}} \Leftrightarrow$ for all $A \subseteq B$ finite there is $g \in \mathcal{S}$ such that $g|_A = f|_A$.

$\overline{\langle \mathcal{S} \rangle}$ denotes the smallest closed clone containing \mathcal{S} .

There is a correspondence between:

- closed subclones of \mathcal{O} ;
- polymorphism clones of relational structures on B .

Minimal clones and operations I

Let $\mathcal{D} \supsetneq \mathcal{C}$ be closed subclones of \mathcal{O} .

Definition 2 (almost minimal)

The k -ary operation $f \in \mathcal{D} \setminus \mathcal{C}$ is **almost minimal above \mathcal{C}** if for each $r < k$,

$$\overline{\langle \mathcal{C} \cup \{f\} \rangle} \cap \mathcal{O}^{(r)} = \mathcal{C} \cap \mathcal{O}^{(r)}.$$

Definition 3 (Minimal operation)

$f \in \mathcal{D} \setminus \mathcal{C}$ is **minimal above \mathcal{C}** if it is almost minimal above \mathcal{C} , and everything new that it generates with \mathcal{C} also generates f with \mathcal{C} :

$$h \in \overline{\langle \mathcal{C} \cup \{f\} \rangle} \setminus \mathcal{C} \Rightarrow f \in \overline{\langle \mathcal{C} \cup \{h\} \rangle}.$$

Minimal clones and operations II

Definition 4

\mathcal{D} is **minimal above** \mathcal{C} if there is no closed clone \mathcal{E} such that $\mathcal{C} \subsetneq \mathcal{E} \subsetneq \mathcal{D}$.

- There is a correspondence between
 - minimal closed clones above \mathcal{C} ;
 - clones $\overline{\langle \mathcal{C} \cup \{f\} \rangle}$ for f minimal above \mathcal{C} .
- For $\mathcal{C} \subsetneq \mathcal{D}$, there is always $f \in \mathcal{D}$ **almost minimal** above \mathcal{C} ;
- For the structures B we are most interested in¹, $\text{Pol}(B)$ always has a **minimal** operation above $\overline{\langle \text{Aut}(B) \rangle}$;

¹i.e. B is finite or ω -categorical in a finite relational language.

Some terminology for operations

We define some operations in virtue of the identities they satisfy:

- **ternary quasi-majority:**

$$m(x, x, y) \approx m(x, y, x) \approx m(y, x, x) \approx m(x, x, x);$$

- **quasi-Malcev:**

$$M(y, y, x) \approx M(x, y, y) \approx M(x, x, x);$$

- For f **idempotent**, i.e. $f(x, \dots, x) \approx x$, remove the 'quasi';
- **ternary minority:**

$$\mathfrak{m}(y, y, x) \approx \mathfrak{m}(y, x, y) \approx \mathfrak{m}(x, y, y) \approx \mathfrak{m}(x, x, x) \approx x;$$

- **semiprojection:** a k -ary f such that there is an $i \in \{1, \dots, k\}$ such that whenever $|\{a_1, \dots, a_k\}| < k$,

$$f(a_1, \dots, a_k) = a_i.$$

Rosenberg's five types theorem

Theorem 5 (Five types theorem, Rosenberg 1986)

Let B be finite and f be minimal above $\langle 1 \rangle$. Then, f is of one of the following five types:

- ① *a unary function;*
- ② *a binary function;*
- ③ *a ternary majority operation;*
- ④ *a ternary minority operation;*
- ⑤ *a k -ary semiprojection for some $k \geq 3$.*

Moreover, the ternary minority corresponds to $x + y + z$ in some Boolean group.

A group is **Boolean** if every element has order 2.

Our project

For $G \curvearrowright B$ a non-trivial group G acting faithfully on B , we want to classify the minimal operations lying above $\overline{\langle G \rangle}$.

STRATEGY: Classify almost minimal operations first.

The classification of almost minimal operations splits into three cases:

- G is not a Boolean group acting freely on B ;
- G is a Boolean group acting freely on B and $\neq \mathbb{Z}_2$;
- \mathbb{Z}_2 acting freely on B .

For minimal operations, the last two cases can be treated together.

Classifying almost minimal operations I

Theorem 6 (Three Types Theorem, Marimon and Pinsker 2024)

Let B be a *finite* set. Let $G \curvearrowright B$ be *such that G is not a Boolean group acting freely on B* . Let f be *almost minimal* above $\overline{\langle G \rangle}$. Then, f is one of:

- ① a unary function;
- ② a binary function;
- ③ a k -ary orbit-semiprojection for $3 \leq k \leq s$, where $s = |\text{Orb}(G)|$.

f is an **orbit-semiprojection** if there is an $i \in \{1, \dots, k\}$ and a unary operation $g \in \overline{G}$ such that **whenever at least two of the a_j lie in the same orbit,**

$$f(a_1, \dots, a_k) = g(a_i).$$

Classifying almost minimal operations II

Theorem 7 (Boolean case, Marimon and Pinsker 2024)

Let $G \curvearrowright B$ be a Boolean group acting freely on B with s -many orbits and $|G| > 2$. Let f be an almost minimal operation above $\langle G \rangle$. Then, f is of one of the following types:

- ① f is unary;
- ② f is binary;
- ③ f is a **ternary twisted minority**;
- ④ f is a k -ary orbit-semiprojection for $3 \leq k \leq s$.

A **twisted minority** is a ternary operation such that for all $\beta \in G$,

$$\mathbf{m}(y, x, \beta x) \approx \mathbf{m}(x, \beta x, y) \approx \mathbf{m}(x, y, \beta x) \approx \mathbf{m}(\beta y, \beta y, \beta y).$$

Classifying almost minimal operations III

Theorem 8 (\mathbb{Z}_2 case, Marimon and Pinsker 2024)

Let \mathbb{Z}_2 act freely on B with s -many orbits. Let f be an almost minimal operation above $\langle \mathbb{Z}_2 \rangle$. Then, f is of one of the following types:

- ① f is unary;
- ② f is a ternary twisted minority;
- ③ f is an **odd majority**;
- ④ f is, up to permuting its variables, an **odd Malcev**;
- ⑤ f is a k -ary orbit-semiprojection for $2 \leq k \leq s$.

An **odd majority** m is a quasi-majority such that for γ the non-identity element in \mathbb{Z}_2 ,

$$m(y, x, \gamma x) \approx m(x, \gamma x, y) \approx m(x, y, \gamma x) \approx m(y, y, y).$$

An **odd Malcev** is a quasi-Malcev such that $M(x, \gamma y, z)$ is an odd majority.

Minimal operations

- **Orbit semiprojections** only generate other orbit semiprojections, so for $|\text{Orb}(G)| \geq 3$, they always exist as minimal;²
- **Odd majorities** and **odd Malcev**s are **NEVER** minimal;
- **Minimal twisted minorities** are particularly well-behaved:
 - $\mathbf{m}(x, y, z) = w$ is a symmetric 4-ary relation;
 - For $\alpha, \beta, \gamma \in G$, $\mathbf{m}(\alpha x, \beta y, \gamma z) \approx \alpha\beta\gamma\mathbf{m}(x, y, z)$;
 - \mathbf{m} induces on the orbits of G a minority of the form $\mathbf{m}^*(x, y, z) = x + y + z$ in a Boolean group;
 - They exist above $\langle G \rangle$ if and only if $|\text{Orb}(G)| = 2^n$ or is infinite.

²At least in the contexts that interest us.

Applications: finding low arity essential polymorphisms I

We want to study $\text{CSP}(B)$ for B countably infinite and $\text{Aut}(B) \curvearrowright B$

oligomorphic: $\text{Aut}(B) \curvearrowright B^n$ has finitely many orbits for all n .

We call these structures **ω -categorical**.

Sufficient to work with a **model complete core**: $\overline{\text{Aut}(B)} = \text{End}(B)$.

It is often helpful to find low arity essential polymorphisms:

Definition 9 (Essentially unary and essential operations)

f is **essentially unary** if there is unary g and $1 \leq i \leq k$ such that

$$f(x_1, \dots, x_k) \approx g(x_i).$$

Otherwise, f is **essential**.

Applications: finding low arity essential polymorphisms II

For B ω -categorical, people developed techniques to find binary essential polymorphisms given an essential one (cf. Bodirsky and Kára 2008, and Mottet and Pinsker 2022).

Question 1 (Question 24 in Bodirsky 2021)

Does an ω -categorical model complete core with an essential polymorphism also have a binary essential polymorphism?

Answer:

- **NO:** whenever $\text{Aut}(B)$ has ≥ 3 orbits, $\overline{\langle \text{Aut}(B) \cup \{f\} \rangle}$ for f an orbit semiprojection is a counterexample;
- However, this is true on ≤ 2 orbits;
- Moreover, this is true when $\text{CSP}(B)$ is not hard...

Applications: finding low arity essential polymorphisms III

Theorem 10 (Marimon and Pinsker 2024)






Suppose B is a finite or ω -categorical model complete core and $\text{Aut}(B) \curvearrowright B$ is not the free action of a Boolean group on B (always the case if B is ω -categorical).

(\star) Suppose $\text{Pol}(B)$ does not have a uniformly continuous $h1$ -homomorphism to Proj (the clone of projections on $\{0, 1\}$);

Then, $\text{Pol}(B)$ contains a binary essential polymorphism.

- $\neg(\star) \Rightarrow \text{CSP}(B)$ is NP hard;
- Proof uses the Three Types Theorem (for almost minimal operations).

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Funding Statement: Funded by the European Union (ERC, POCOCOP, 101071674). Views and opinions expressed are however those of the authors only and do not necessarily reflect those of the European Union or the European Research Council Executive Agency. Neither the European Union nor the granting authority can be held responsible for them.