

Clones over finite sets and Minor Conditions

Albert Vucaj

TU Wien



ERC SYNERGY GRANT
POCOP; 101071674

Linz , 17/10/2023

Clones

Def: A clone over a (finite) set A is a set of operations \mathcal{A} over A s.t.

- \mathcal{A} contains all projections;
- \mathcal{A} is closed under composition.

Let F be a set of operations, we denote by $\langle F \rangle$ the clone generated by F .

Clones

Def: A clone over a (finite) set A is a set of operations \mathcal{A} over A s.t.

- \mathcal{A} contains all projections;
- \mathcal{A} is closed under composition.

Let F be a set of operations, we denote by $\langle F \rangle$ the clone generated by F .

<u>Example:</u>	$P_2 := \langle \emptyset \rangle$	all projections over $\{0,1\}$;
	$I_2 := \langle \wedge; z \mapsto y \mapsto z \rangle$	all idempotent operations over $\{0,1\}$;
	$1 := \langle 1 \rangle$	the clone generated by the constant operation 1.

Clones

Def: A clone over a (finite) set A is a set of operations \mathcal{A} over A s.t.

- \mathcal{A} contains all projections;
- \mathcal{A} is closed under composition.

Let F be a set of operations, we denote by $\langle F \rangle$ the clone generated by F .

Example: $P_2 := \langle \emptyset \rangle$ all projections over $\{0,1\}$;
 $I_2 := \langle \wedge; z \rightarrow y \rightarrow z \rangle$ all idempotent operations over $\{0,1\}$;
 $1 := \langle 1 \rangle$ the clone generated by the constant operation 1.

Def: An operation $f: A^n \rightarrow A$ preserves a relation $R \subseteq A^k$ if $\forall a_1, \dots, a_n \in R$, $\begin{pmatrix} f(a_1^1, \dots, a_n^1) \\ \vdots \\ f(a_1^K, \dots, a_n^K) \end{pmatrix} \in R$.

Clones

Def: A clone over a (finite) set A is a set of operations \mathcal{A} over A s.t.

- \mathcal{A} contains all projections;
- \mathcal{A} is closed under composition.

Let F be a set of operations, we denote by $\langle F \rangle$ the clone generated by F .

Example: $P_2 := \langle \emptyset \rangle$ all projections over $\{0,1\}$;

$I_2 := \langle \wedge; z \rightarrow y \rightarrow z \rangle$ all idempotent operations over $\{0,1\}$;

$1 := \langle 1 \rangle$ the clone generated by the constant operation 1.

Def: An operation $f: A^n \rightarrow A$ preserves a relation $R \subseteq A^k$ if $\forall a_1, \dots, a_n \in R, \begin{pmatrix} f(a_1^1, \dots, a_n^1) \\ \vdots \\ f(a_1^K, \dots, a_n^K) \end{pmatrix} \in R$.

• An operation f is a polymorphism of $\mathbb{A} = (A; \Gamma)$ if f preserves R , for every $R \in \Gamma$.

• By $\text{Pol}(A)$ we denote the set of all polymorphisms of \mathbb{A} . (it is a clone)

Example: $C_p := \text{Pol}(\{0, 1, \dots, p-1\}; \{(0,1), (1,2), \dots, (p-1,0)\})$

$B_2 := \text{Pol}(\{0, 1\}; \{(0,1), (1,0), (1,1)\}, \{0\}, \{1\})$



A Galois Connection for Clones

On the relational side \rightsquigarrow denote by $\text{Inv}(F)$ the set $\{R \mid f \text{ preserves } R, \forall f \in F\}$.

A Galois Connection for Clones

On the relational side \rightsquigarrow denote by $\text{Inv}(F)$ the set $\{R \mid f \text{ preserves } R, \forall f \in F\}$.

- Def:
- A relation $S \subseteq A^n$ is pp-definable from Γ (set of relations over A) if it can be defined by a formula that uses relations from Γ , \exists , \wedge , and $=$.
 - A relational clone over A is a set of relations over A closed under pp-definable relations (and containing \emptyset).

A Galois Connection for Clones

On the relational side \rightsquigarrow denote by $\text{Inv}(F)$ the set $\{R \mid f \text{ preserves } R, \forall f \in F\}$.

- Def:
- A relation $S \subseteq A^n$ is pp-definable from Γ (set of relations over A) if it can be defined by a formula that uses relations from Γ , \exists , \wedge , and $=$.
 - A relational clone over A is a set of relations over A closed under pp-definable relations (and containing \emptyset).

: $\text{Inv}(F)$ is a relational clone.

A Galois Connection for Clones

On the relational side \rightsquigarrow denote by $\text{Inv}(F)$ the set $\{R \mid f \text{ preserves } R, \forall f \in F\}$.

- Def:
- A relation $S \subseteq A^n$ is pp-definable from Γ (set of relations over A) if it can be defined by a formula that uses relations from Γ , \exists, \wedge , and $=$.
 - A relational clone over A is a set of relations over A closed under pp-definable relations (and containing \emptyset).

: $\text{Inv}(F)$ is a relational clone.

Thm [Bodnaruk, Kaluznin, Kotov, Roman; Geiger]: F : set of oper. over some finite universe; Γ : set of relations over some finite universe.

$$\text{Then } \text{Pol}(\text{Inv}(F)) = \langle F \rangle \quad \text{and} \quad \text{Inv}(\text{Pol}(\Gamma)) = [\Gamma] \quad \begin{matrix} \text{relational clone} \\ \text{generated by } R \end{matrix}$$

A Galois Connection for Clones

On the relational side \rightsquigarrow denote by $\text{Inv}(F)$ the set $\{R \mid f \text{ preserves } R, \forall f \in F\}$.

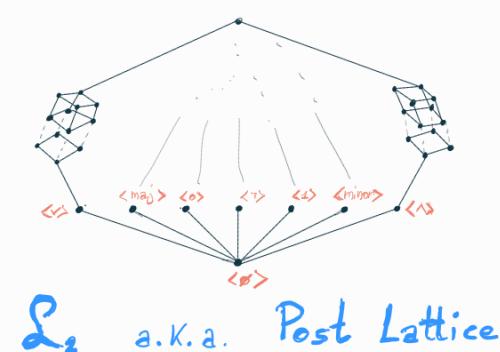
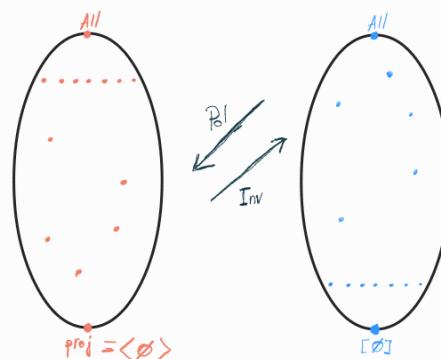
- Def:
- A relation $S \subseteq A^n$ is pp-definable from Γ (set of relations over A) if it can be defined by a formula that uses relations from Γ , \exists, \wedge , and $=$.
 - A **relational clone** over A is a set of relations over A closed under pp-definable relations (and containing \emptyset).

$\text{Inv}(F)$ is a relational clone.

Thm [Bodnaruk, Kaluznin, Kotor, Roman; Geiger]: F : set of oper. over some finite universe; Γ : set of relations over some finite universe.

Then $\text{Pol}(\text{Inv}(F)) = \langle F \rangle$ and $\text{Inv}(\text{Pol}(\Gamma)) = [\Gamma]$ relational clone generated by R

Thm: Let A be an n -element set. The set of all clones over A forms a lattice \mathcal{L}_n under inclusion.



Clones over $\{0, 1, 2\}$

Thm [Yanov, Muchnik]: Let A be a finite set s.t. $|A| \geq 3$. There are 2^ω many clones over A .

Clones over $\{0,1,2\}$

Thm [Yanov, Muchnik]: Let A be a finite set s.t. $|A| \geq 3$. There are 2^ω many clones over A .

Subsequent research focused on describing some parts of the lattice \mathcal{L}_3 :

- Maximal clones [Jablonskij, '54]
- Minimal clones [Csákány, '83]

Clones over $\{0,1,2\}$

Thm [Yanov, Muchnik]: Let A be a finite set s.t. $|A| \geq 3$. There are 2^ω many clones over A .

Subsequent research focused on describing some parts of the lattice \mathcal{L}_3 :

- Maximal clones [Jablonskij, '54]
- Minimal clones [Csákány, '83]

All maximal clones – except the clone of all linear functions – contain a continuum of subclones.

Clones over $\{0,1,2\}$

Thm [Yanov, Muchnik]: Let A be a finite set s.t. $|A| \geq 3$. There are 2^ω many clones over A .

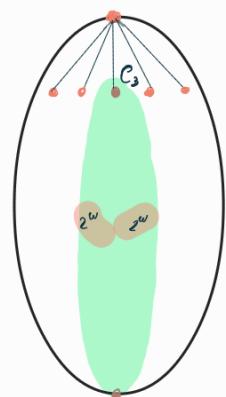
Subsequent research focused on describing some parts of the lattice \mathcal{L}_3 :

- Maximal clones [Jablonskij, '54]
- Minimal clones [GáKány, '83]

All maximal clones - except the clone of all linear functions - contain a continuum of subclones.

D. Zhuk completely described the lattice of subclones of $C_3 := \text{Pol}(\cdot \circ \cdot)$ [2015].

the clone of all self dual operations over $\{0,1,2\}$.



Clones over $\{0,1,2\}$

Thm [Yanov, Muchnik]: Let A be a finite set s.t. $|A| \geq 3$. There are 2^ω many clones over A .

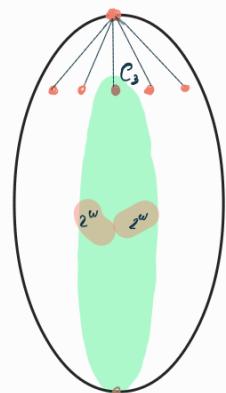
Subsequent research focused on describing some parts of the lattice \mathcal{L}_3 :

- Maximal clones [Jablonskij, '54]
- Minimal clones [GáKáng, '83]

All maximal clones - except the clone of all linear functions - contain a continuum of subclones.

D. Zhuk completely described the lattice of subclones of $C_3 := \text{Pol}(\cdot \circ \cdot)$ [2015].

the clone of all self dual operations over $\{0,1,2\}$.



💡: C_3 is the only maximal clone over $\{0,1,2\}$, containing 2^ω many subclones whose lattice of subclones has been completely described.

A New Order

Def: A minor of an operation f is any operation obtained from f by

identifying variables
permuting variables
adding new variables

A New Order

Def: A **minor** of an operation f is any operation obtained from f by

identifying variables
permuting variables
adding new variables

Formally, let f be an n -ary operation and $\pi: \{1, \dots, n\} \rightarrow \{1, \dots, r\}$.

We write f_π to denote the r -ary operation $f_\pi(x_1, \dots, x_r) := f(x_{\pi(1)}, \dots, x_{\pi(n)})$.

A New Order

Def: A **minor** of an operation f is any operation obtained from f by

identifying variables
permuting variables
adding new variables

Formally, let f be an n -ary operation and $\pi: \{1, \dots, n\} \rightarrow \{1, \dots, r\}$.

We write f_π to denote the r -ary operation $f_\pi(x_1, \dots, x_r) := f(x_{\pi(1)}, \dots, x_{\pi(n)})$.

Def: A **minor preserving** map is a map $\Xi: A \rightarrow B$ ($A \leq_m B$) such that

- Ξ preserves arities;
- $\Xi(f_\pi) = \Xi(f)_\pi$, for any n -ary $f \in A$ and $\pi: \{1, \dots, n\} \rightarrow \{1, \dots, r\}$.

$$\begin{array}{ccc} f & \xrightarrow{\Xi} & \Xi(f) \\ \downarrow \pi & & \downarrow \pi \\ f_\pi & \xrightarrow{\Xi} & \Xi(f_\pi) = \Xi(f)_\pi \end{array}$$

A New Order

Def: A **minor** of an operation f is any operation obtained from f by

identifying variables
permuting variables
adding new variables

Formally, let f be an n -ary operation and $\pi: \{1, \dots, n\} \rightarrow \{1, \dots, r\}$.

We write f_π to denote the r -ary operation $f_\pi(x_1, \dots, x_r) := f(x_{\pi(1)}, \dots, x_{\pi(n)})$.

Def: A **minor preserving** map is a map $\Xi: A \rightarrow B$ ($A \leq_m B$) such that

- Ξ preserves arities;
- $\Xi(f_\pi) = \Xi(f)_\pi$, for any n -ary $f \in A$ and $\pi: \{1, \dots, n\} \rightarrow \{1, \dots, r\}$.

$$\begin{array}{ccc} f & \xrightarrow{\Xi} & \Xi(f) \\ \downarrow \pi & & \downarrow \pi \\ f_\pi & \xrightarrow{\Xi} & \Xi(f_\pi) = \Xi(f)_\pi \end{array}$$

On the relational side:

Def: We say that B is a **pp-power** of A if B is isomorphic to a structure P s.t.

- the domain of P is A^n , for some $n \geq 1$;
- all the relations of P are pp-definable from A .

A New Order

Def: A **minor** of an operation f is any operation obtained from f by

identifying variables
permuting variables
adding new variables

Formally, let f be an n -ary operation and $\pi: \{1, \dots, n\} \rightarrow \{1, \dots, r\}$.

We write f_π to denote the r -ary operation $f_\pi(x_1, \dots, x_r) := f(x_{\pi(1)}, \dots, x_{\pi(n)})$.

Def: A **minor preserving** map is a map $\Xi: A \rightarrow B$ ($A \leq_m B$) such that

- Ξ preserves arities;
- $\Xi(f_\pi) = \Xi(f)_\pi$, for any n -ary $f \in A$ and $\pi: \{1, \dots, n\} \rightarrow \{1, \dots, r\}$.

$$\begin{array}{ccc} f & \xrightarrow{\Xi} & \Xi(f) \\ \downarrow \pi & & \downarrow \pi \\ f_\pi & \xrightarrow{\Xi} & \Xi(f_\pi) = \Xi(f)_\pi \end{array}$$

On the relational side:

Def: We say that B is a **pp-power** of A if B is isomorphic to a structure P s.t.

- the domain of P is A^n , for some $n \geq 1$;
- all the relations of P are pp-definable from A .

Def: A **pp-constructs** B if B is hom. equivalent to a pp-power of A .

CSP in a Nutshell

Def: A, B : τ -structures (τ : finite relational signature).

A **homomorphism** from A to B ($A \rightarrow B$) is a map $h: A \rightarrow B$ st., for every $R \in \tau$,

$$(a_1, \dots, a_n) \in R^A \implies (h(a_1), \dots, h(a_n)) \in R^B$$

CSP in a Nutshell

Def: A, B : τ -structures (τ : finite relational signature).

A **homomorphism** from A to B ($A \rightarrow B$) is a map $h: A \rightarrow B$ st., for every $R \in \tau$,

$$(a_1, \dots, a_n) \in R^A \implies (h(a_1), \dots, h(a_n)) \in R^B$$

Def: For a fixed relational τ -structure A , **CSP(A)** is the membership problem of

$$\{ S \mid S \text{ is a } \tau\text{-structure and } S \rightarrow A \}$$

CSP in a Nutshell

Def: A, B : τ -structures (τ : finite relational signature).

A **homomorphism** from A to B ($A \rightarrow B$) is a map $h: A \rightarrow B$ st., for every $R \in \tau$,

$$(a_1, \dots, a_n) \in R^A \implies (h(a_1), \dots, h(a_n)) \in R^B$$

Def: For a fixed relational τ -structure A , **CSP(A)** is the membership problem of

$$\{S \mid S \text{ is a } \tau\text{-structure and } S \rightarrow A\}$$

Example: • $CSP(K_3)$ is equivalent to the 3-colorability problem for graphs.



• $CSP(\underbrace{\{0, \dots, p-1\}}_{E_p}; \text{all affine subspaces } R_{abcd}) \sim$ solving systems of lin. eq. over \mathbb{Z}_p

$$R_{abcd} := \{(x, y, z) \in E_p^3 \mid ax + by + cz = d\}$$

Minor Conditions

Thm [G '68; BKRR '69]:

$$A \text{ pp-defines } B \Leftrightarrow \text{Pol}(A) \subseteq \text{Pol}(B)$$

Minor Conditions

Thm [G. '68; BKKR '69] :

$$A \text{ pp-defines } B \Leftrightarrow \text{Pol}(A) \subseteq \text{Pol}(B)$$

Thm [Barto, Oprea, Pinski '15] :

$$A \text{ pp-constructs } B \Leftrightarrow \text{Pol}(A) \leq_m \text{Pol}(B)$$

Minor Conditions

Thm [G. '68; BKKR '69] :

$$A \text{ pp-defines } B \Leftrightarrow \text{Pol}(A) \subseteq \text{Pol}(B)$$

Thm [Barto, Oprea, Pinski '15] :

$$A \text{ pp-constructs } B \Leftrightarrow \text{Pol}(A) \leq_m \text{Pol}(B)$$

Thm [Barto, Oprea, Pinski '15] : If $\text{Pol}(A) \leq_m \text{Pol}(B)$, then $CSP(B)$ reduces to $CSP(A)$.

Minor Conditions

Thm [G. '68; BKKR '69] :

$$A \text{ pp-defines } B \Leftrightarrow \text{Pol}(A) \subseteq \text{Pol}(B)$$

Thm [Barto, Opršal, Pinski '15] :

$$A \text{ pp-constructs } B \Leftrightarrow \text{Pol}(A) \leq_m \text{Pol}(B)$$

Thm [Barto, Opršal, Pinski '15] : If $\text{Pol}(A) \leq_m \text{Pol}(B)$, then $\text{CSP}(B)$ reduces to $\text{CSP}(A)$.

Def : A **minor identity** is an abstract expression of the form

$$\forall x_1, \dots, x_n, y_1, \dots, y_m \quad (\underbrace{f(x_1, \dots, x_n)}_{\text{operation symbols}} = \underbrace{g(y_1, \dots, y_m)}_{\text{not necessarily distinct}})$$

For short:

$$f(x_1, \dots, x_n) \approx g(y_1, \dots, y_m)$$
or
$$f \approx g$$

Minor Conditions

Thm [G. '68; BKKR '69]:

$$A \text{ pp-defines } B \Leftrightarrow \text{Pol}(A) \subseteq \text{Pol}(B)$$

Thm [Barto, Oprea, Pinski '15]:

$$A \text{ pp-constructs } B \Leftrightarrow \text{Pol}(A) \leq_m \text{Pol}(B)$$

Thm [Barto, Oprea, Pinski '15]: If $\text{Pol}(A) \leq_m \text{Pol}(B)$, then $\text{CSP}(B)$ reduces to $\text{CSP}(A)$.

Def: A **minor identity** is an abstract expression of the form

$$\forall x_1, \dots, x_n, y_1, \dots, y_m \quad (\underbrace{f(x_1, \dots, x_n)}_{\text{operation symbols}} = \underbrace{g(y_1, \dots, y_m)}_{\text{not necessarily distinct}})$$

For short:
 $f(x_1, \dots, x_n) \approx g(y_1, \dots, y_m)$
 or
 $f \approx g$

- A **minor condition** is a finite set of minor identities.

Example: $m(x, x, y) \approx m(y, x, x) \approx m(y, y, y)$
 QUASI MAJ'CEV

$c(x_1, \dots, x_p) \approx c(x_2, \dots, x_p, x_1)$
 p-CYCLIC

- $m(x, x, y) \approx m(x, y, x) \approx m(y, x, x) \approx m(x, x, x)$
 QUASI MAJORITY
- $f(\underbrace{x_1, \dots, x_n}_X) \approx f(\underbrace{y_1, \dots, y_n}_Y)$, whenever $X = Y$.
 TOL. SYMMETRIC

Algebra meets CSP

- Def: We say the a set of operations F satisfies a minor condition Σ if there is a map $\tilde{\Sigma}$ assigning to each operation symbol occurring in F an operation in F of the same variety s.t. if $(p \approx q) \in \Sigma$, then $\tilde{\Sigma}(p) = \tilde{\Sigma}(q)$. ($F \models \Sigma$)
- A minor condition Σ is trivial if $P_2 \models \Sigma$.

Algebra meets CSP

Def: We say that a set of operations F satisfies a minor condition Σ if there is a map $\tilde{\Sigma}$ assigning to each operation symbol occurring in F an operation in F of the same arity s.t. if $(p \approx q) \in \Sigma$, then $\tilde{\Sigma}(p) = \tilde{\Sigma}(q)$. ($F \models \Sigma$)

- A minor condition Σ is trivial if $P_2 \models \Sigma$.

[Barto, Oporšal, Pinsker]: The complexity of $CSP(A)$ only depends on the set of minor conditions satisfied in $Pol(A)$.

Algebra meets CSP

Def: We say the a set of operations F satisfies a minor condition Σ if there is a map \gtrsim assigning to each operation symbol occurring in F an operation in F of the same arity s.t. if $(p \approx q) \in \Sigma$, then $\gtrsim(p) = \gtrsim(q)$. ($F \models \Sigma$)

- A minor condition Σ is trivial if $P_2 \models \Sigma$.

[Barto, Oprea, Pinsker]: The complexity of $CSP(A)$ only depends on the set of minor conditions satisfied in $\text{Pol}(A)$.

Thm [BOP '15]: Let A and B be finite relational structures. TFAE:

- (1) There exist a minor-preserving map from $\text{Pol}(A)$ to $\text{Pol}(B)$ ($\text{Pol}(A) \leq_m \text{Pol}(B)$);
- (2) A pp-constructs B ($A \leq_{\text{con}} B$);
- (3) If $\text{Pol}(A) \models \Sigma$, then $\text{Pol}(B) \models \Sigma$.

Algebra meets CSP

Great achievement: CSP Dichotomy Theorem (Bulatov '17; Zhuk '17)

Algebra meets CSP

Great achievement: CSP Dichotomy Theorem (Bulatov '17; Zhuk '17)

- ▶ positive answer to the Feder-Vardi conjecture open since 1998;

Algebra meets CSP

Great achievement: CSP Dichotomy Theorem (Bulatov '17; Zhuk '17)

- ▶ positive answer to the Feder-Vardi conjecture open since 1998;
- ▶ new algebraic theories for finite algebras
 - absorption (Barto, Kozik)
 - Bulatov-edges (Bulatov)
 - Strong subalgebras
 - ...

Algebra meets CSP

Great achievement: CSP Dichotomy Theorem (Bulatov '17; Zhuk '17)

- ▶ positive answer to the Feder-Vardi conjecture open since 1998;
- ▶ new algebraic theories for finite algebras
 - absorption (Barto, Kozik)
 - Bulatov-edges (Bulatov)
 - Strong subalgebras
 - ...

Thm [Siggers 2010]: Let C be a clone over a finite set. TFAE:

- (1) C satisfies a non-trivial minor condition;
- (2) $C \models s(x, y, z, z, y, z) \approx s(y, x, x, z, z, y)$.

Algebra meets CSP

Great achievement: CSP Dichotomy Theorem (Bulatov '17; Zhuk '17)

- ▶ positive answer to the Feder-Vardi conjecture open since 1998;
- ▶ new algebraic theories for finite algebras
 - absorption (Barto, Kozik)
 - Bulatov-edges (Bulatov)
 - Strong subalgebras
 - ...

Thm [Siggers 2010]: Let C be a clone over a finite set. TFAE:

- (1) C satisfies a non-trivial minor condition;
- (2) $C \models s(x, y, z, x, y, z) \approx s(y, x, x, z, z, y)$.

Thm [Bulatov '17; Zhuk '17]: Let $A = \text{Pol}(A)$ be a clone over a finite set.

If $A \models s(x, y, z, x, y, z) \approx s(y, x, x, z, z, y)$, then $\text{CSP}(A)$ is in P.

Otherwise, $\text{CSP}(A)$ is NP-complete

The pp-constructability poset

We write $C \equiv_m D$ iff $C \leq_m D$ and $D \leq_m C$. it is reflexive and transitive!

The pp-constructability poset

We write $C \equiv_m D$ iff $C \leq_m D$ and $D \leq_m C$. it is reflexive and transitive!

We denote by \bar{C} the \equiv_m -class of C . Also, we write $\bar{C} \leq_m \bar{D}$ iff $C \leq_m D$.

$$P_{\text{fin}} := (\{\bar{C} \mid C \text{ is a c.c. over some finite set}\}; \leq_m)$$

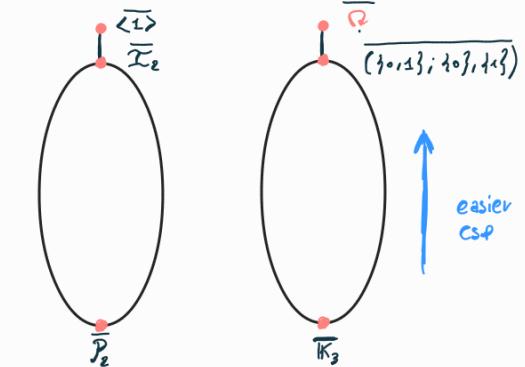
The pp-constructability poset

We write $C \equiv_m D$ iff $C \leq_m D$ and $D \leq_m C$. it is reflexive and transitive!

We denote by \bar{C} the \equiv_m -class of C . Also, we write $\bar{C} \leq_m \bar{D}$ iff $C \leq_m D$.

$$\mathcal{P}_{\text{fin}} := (\{\bar{C} \mid C \text{ is a c.c. over some finite set}\}; \leq_m)$$

$$\mathcal{P}_n := (\{\bar{C} \mid C \text{ is a c.c. over } \{0, 1, \dots, n-1\}\}; \leq_m)$$



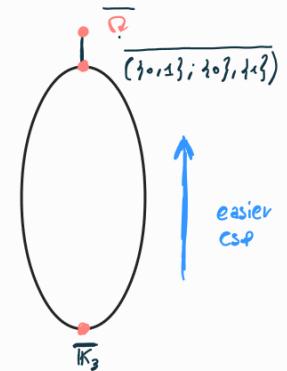
The pp-constructability poset

We write $C \equiv_m D$ iff $C \leq_m D$ and $D \leq_m C$. it is reflexive and transitive!

We denote by \bar{C} the \equiv_m -class of C . Also, we write $\bar{C} \leq_m \bar{D}$ iff $C \leq_m D$.

$$\mathcal{P}_{\text{fin}} := (\{\bar{C} \mid C \text{ is a clone over some finite set}\}; \leq_m)$$

$$\mathcal{P}_n := (\{\bar{C} \mid C \text{ is a clone over } \{0, 1, \dots, n-1\}\}; \leq_m)$$



Some results: ▶ Complete description of \mathcal{P}_2 [Bodirsky, V.]



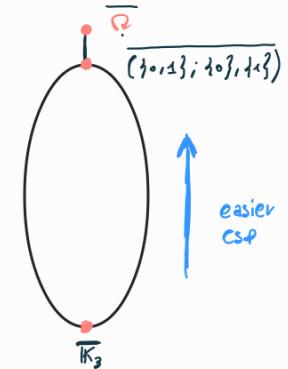
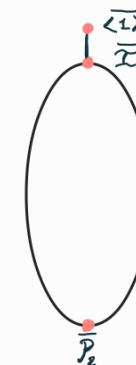
The pp-constructability poset

We write $C \equiv_m D$ iff $C \leq_m D$ and $D \leq_m C$. it is reflexive and transitive!

We denote by \bar{C} the \equiv_m -class of C . Also, we write $\bar{C} \leq_m \bar{D}$ iff $C \leq_m D$.

$$\mathcal{P}_{\text{fin}} := (\{\bar{C} \mid C \text{ is a dome over some finite set}\}; \leq_m)$$

$$\mathcal{P}_n := (\{\bar{C} \mid C \text{ is a dome over } \{0, 1, \dots, n-1\}\}; \leq_m)$$



- Some results:
- Complete description of \mathcal{P}_2 [Bodirsky, V.]
 - Complete description of \mathcal{P}_{SD} (pol. domes of smooth digraphs) [Bodirsky, Stark, V.]



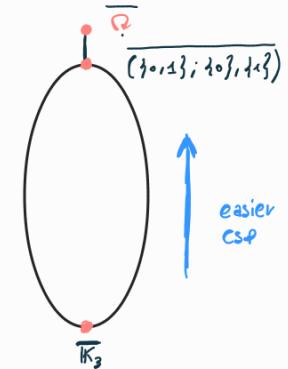
The pp-constructability poset

We write $C \equiv_m D$ iff $C \leq_m D$ and $D \leq_m C$. it is reflexive and transitive!

We denote by \bar{C} the \equiv_m -class of C . Also, we write $\bar{C} \leq_m \bar{D}$ iff $C \leq_m D$.

$$\mathcal{P}_{\text{fin}} := (\{\bar{C} \mid C \text{ is a clone over some finite set}\}; \leq_m)$$

$$\mathcal{P}_n := (\{\bar{C} \mid C \text{ is a clone over } \{0, 1, \dots, n-1\}\}; \leq_m)$$



- Some results:
- ▶ Complete description of \mathcal{P}_2 [Bodirsky, V.]
 - ▶ Complete description of \mathcal{P}_{SD} (pol. clones of smooth digraphs) [Bodirsky, Stark, V.]
 - ▶ Complete description of the lattice of clones of self-dual clones w.r.t. \leq_m [Bodirsky, Zhu, V.]
- remarkably
countably infinite!



Submaximal elements of \mathcal{P}_3

Want: Show that $\overline{\text{Pol}(\uparrow\downarrow)}$, $\overline{\text{Pol}(\uparrow\downarrow\downarrow)}$, $\overline{\text{Pol}(\{1_0, 1\}; \{0_1, 1_0, 1_1, 1_2\}, 1_0, 1_1)}$ are the only submax. el. of \mathcal{P}_3 .

$$\text{Pol}(\uparrow\downarrow\downarrow) = \overline{\mathcal{C}_2}$$

$$\overline{\mathcal{C}_3}$$

$$\overline{\mathcal{B}_2} = \text{Pol}(\uparrow\downarrow\downarrow)$$

Submaximal elements of \mathcal{P}_3

Want: Show that $\overline{\text{Pol}(\uparrow\downarrow)}$, $\overline{\text{Pol}(\uparrow\downarrow\downarrow)}$, $\overline{\text{Pol}(\{1_0, 1\}; \{0_1, 1_1, 1_2\}, 1_0, 1_1)}$ are the only submax. el. of \mathcal{P}_3 .

$$\text{Pol}(\uparrow\downarrow\downarrow) = \overline{\mathcal{C}_2}$$

$$\overline{\mathcal{C}_3}$$

$$\overline{\mathcal{B}_2} = \text{Pol}(\uparrow\downarrow\downarrow)$$

What we really want: Show that if A is a dom over $\{0, 1, 2\}$ such that $A \not\leq_m \mathcal{C}_2$, $A \not\leq_m \mathcal{C}_3$, and $A \not\leq_m \mathcal{B}_2$, \star , then $\mathcal{I}_2 \leq_m A$.

Submaximal elements of \mathcal{P}_3

Want: Show that $\overline{\text{Pol}(\text{↑↓})}$, $\overline{\text{Pol}(\text{↔})}$, $\overline{\text{Pol}(\{1_0, 1\}; \{0_1, 1\}, \{1_0, 1_1\}, \{1_0\}, \{1_1\})}$ are the only submax. el. of \mathcal{P}_3 .

$\text{Pol}(\text{↑↓ } \cdot) = \overline{\mathcal{C}_2}$ $\text{Pol}(\text{↔ }) = \overline{\mathcal{C}_3}$ $\overline{\mathcal{B}_2} = \text{Pol}(\text{↑↓})$

What we really want: Show that if A is a dom over $\{0, 1, 2\}$ such that $A \not\leq_m \mathcal{C}_2$, $A \not\leq_m \mathcal{C}_3$, and $A \not\leq_m \mathcal{B}_2$, \star , then $\mathcal{I}_2 \leq_m A$.

Key observations:

- $\mathcal{I}_2 = \langle \wedge, x +_2 y +_2 z \rangle$
- $\forall n \geq 2, \wedge(x_1, \wedge(x_2, \wedge(x_3, \dots \wedge(x_{n-1}, x_n) \dots)))$ is a totally symm. op. of arity n .

Submaximal elements of \mathcal{P}_3

Want: Show that $\overline{\text{Pol}(\text{↑↓})}$, $\overline{\text{Pol}(\text{↔})}$, $\overline{\text{Pol}(\{1_0, 1\}; \{0_1, 1\}, \{1_0, 1_1\}, \{1_0\}, \{1_1\})}$ are the only submax. el. of \mathcal{P}_3 .

$\text{Pol}(\text{↑↓ } \cdot) = \overline{\mathcal{C}_2}$ $\text{Pol}(\text{↔ }) = \overline{\mathcal{C}_3}$ $\overline{\mathcal{B}_2} = \text{Pol}(\text{↑↓})$

What we really want: Show that if A is a dom over $\{0, 1, 2\}$ such that $A \not\leq_m \mathcal{C}_2$, $A \not\leq_m \mathcal{C}_3$, and $A \not\leq_m \mathcal{B}_2$, \star , then $\mathcal{I}_2 \leq_m A$.

- Key observations:
- $\mathcal{I}_2 = \langle 1, x +_2 y +_2 z \rangle$
 - $\forall n \geq 2$, $\wedge(x_1, \wedge(x_2, \wedge(x_3, \dots \wedge(x_{n-1}, x_n) \dots)))$ is a totally symm. op. of arity n .
 - every term operation in \mathcal{I}_2 can be expressed as a sum of l many monomials, for some odd l ; the use of constants is not allowed.

Submaximal elements of \mathcal{P}_3

Want: Show that $\overline{\text{Pol}(\text{↑↓})}$, $\overline{\text{Pol}(\text{↔})}$, $\overline{\text{Pol}(\{1_0, 1\}; \{0_1, 1\}, \{1_0, 1_1\}, \{1_0\}, \{1_1\})}$ are the only submax. el. of \mathcal{P}_3 .

$$\text{Pol}(\text{↑↓ } \cdot) = \overline{\mathcal{C}_2} \quad \text{Pol}(\text{↔ } \cdot) = \overline{\mathcal{C}_3} \quad \overline{\mathcal{B}_2} = \text{Pol}(\text{↑↓})$$

What we really want: Show that if A is a dme over $\{0, 1, 2\}$ such that $A \not\leq_m \mathcal{C}_2$, $A \not\leq_m \mathcal{C}_3$, and $A \not\leq_m \mathcal{B}_2$, \circledast , then $\mathcal{I}_2 \leq_m A$.

- Key observations:
- $\mathcal{I}_2 = \langle 1, x +_2 y +_2 z \rangle$
 - $\forall n \geq 2$, $\wedge(x_1, \wedge(x_2, \wedge(x_3, \dots \wedge(x_{n-1}, x_n) \dots)))$ is a totally symm. op. of arity n .
 - every term operation in \mathcal{I}_2 can be expressed as a sum of l many monomials, for some odd l ; the use of constants is not allowed.

→ show that if A is a dme satisfying \circledast , then A has

- a tot. symm. op. of arity n , $\forall n \geq 2$;
- a operation that "simulates the odd sum modulo 2".

Submaximal elements of \mathbb{P}_3

First step:

Thm: Let A be a finite structure. Then for every prime p :

A pp-contains $C_p \iff \text{Pol}(A)$ does not satisfy $c(x_1, \dots, x_p) \approx c(x_2, \dots, x_p, x_1)$

Σ_p

Submaximal elements of \mathbb{P}_3

First step:

Thm: Let A be a finite structure. Then for every prime p :

A pp-contains $\mathbb{C}_p \iff \text{Pol}(A)$ does not satisfy $c(x_1, \dots, x_p) \approx c(x_2, \dots, x_p, x_1)$ Σ_p

proof (sketch): (\Rightarrow) Suppose $A \leq_{\text{con}} \mathbb{C}_p$, for some prime p .

If $\text{Pol}(A) \models \Sigma_p$ $\xrightarrow{\text{BOP}}$ $\text{Pol}(\mathbb{C}_p) \models \Sigma_p$ i.e.

\Rightarrow we get a pair (c, c) in \mathbb{C}_p

$$\begin{array}{l} f(0, \dots, p-1) = c \\ \downarrow \quad \downarrow \\ f(1, \dots, 0) = c \end{array}$$

Submaximal elements of \mathbb{P}_3

First step:

Thm: Let A be a finite structure. Then for every prime p :

A pp-contains $C_p \iff \text{Pol}(A)$ does not satisfy $c(x_1, \dots, x_p) \approx c(x_2, \dots, x_p, x_1)$ \sum_p

proof (sketch): (\Rightarrow) Suppose $A \leq_{\text{con}} C_p$, for some prime p .

If $\text{Pol}(A) \models \sum_p$ $\xrightarrow{\text{BOP}}$ $\text{Pol}(C_p) \models \sum_p$ i.e.

\Rightarrow we get a pair (c, c) in C_p

$$\begin{array}{l} f(0, \dots, p-1) = c \\ \downarrow \quad \downarrow \\ f(1, \dots, 0) = c \end{array}$$

(\Leftarrow) Suppose $\text{Pol}(A) \not\models \sum_p$. Want to build a pp-power of A , which is hom. eq. to C_p .

- Domain of P , $P = \{f \mid f: A^P \rightarrow A\}$

$$P = (P; R^P)$$

- Relation R^P is the following binary relation:

it consists of all pairs (f, g) s.t. $f, g \in \text{Pol}(A)$ and

$$\begin{array}{l} f(x_1, \dots, x_p) = \\ g(x_2, \dots, x_p, x_1) \end{array}$$

Submaximal elements of \mathbb{P}_3

First step:

Thm: Let A be a finite structure. Then for every prime p :

A pp-contracts $C_p \iff \text{Pol}(A)$ does not satisfy $c(x_1, \dots, x_p) \approx c(x_2, \dots, x_p, x_1)$ \sum_p

proof (sketch): (\Rightarrow) Suppose $A \leq_{\text{con}} C_p$, for some prime p .

If $\text{Pol}(A) \models \sum_p$ $\xrightarrow{\text{BOP}}$ $\text{Pol}(C_p) \models \sum_p$ i.e.

\Rightarrow we get a pair (c, c) in C_p

$$\begin{array}{l} f(0, \dots, p-1) = c \\ \downarrow \quad \downarrow \\ f(1, \dots, 0) = c \end{array}$$

(\Leftarrow) Suppose $\text{Pol}(A) \not\models \sum_p$. Want to build a pp-power of A , which is hom. eq. to C_p .

- Domain of P , $P = \{f \mid f: A^p \rightarrow A\}$

- Relation R^P is the following binary relation:

it consists of all pairs (f, g) s.t. $f, g \in \text{Pol}(A)$ and

$$\begin{array}{l} f(x_1, \dots, x_p) = \\ g(x_2, \dots, x_p, x_1) \end{array}$$

- Partition $\text{Pol}_p(A)$ into eq. classes s.t.

$f_1 \sim f_2$ if f_1 can be obtained by a cyclic shift of the variables of f_2 .

(to be continued) \downarrow

Submaximal elements of \mathbb{P}

proof: • Choose one representative from every \sim -class \rightarrow denote this set of operations F_0 .

Submaximal elements of \mathbb{P}_3

proof: • Choose one representative from every \sim -class \rightarrow denote this set of operations F_0 .

• Define $F_i = \{g \mid \exists f \in F_0 : g(x_1, \dots, x_p) = f(x_i, \dots, x_p, x_{i+1}, \dots, x_{i-1})\}$.

Submaximal elements of \mathbb{P}_3

proof: • Choose one representative from every \sim -class \rightarrow denote this set of operations F_0 .

- Define $F_i = \{g \mid \exists f \in F_0 : g(x_1, \dots, x_p) = f(x_i, \dots, x_p, x_{i+1}, \dots, x_{i-1})\}$.
- $F_i \cap F_j = \emptyset$, whenever $i \neq j$.

Submaximal elements of \mathbb{P}

proof: • Choose one representative from every \sim -class \rightarrow denote this set of operations F_0 .

- Define $F_i = \{g \mid \exists f \in F_0 : g(x_1, \dots, x_p) = f(x_i, \dots, x_p, x_{i+1}, \dots, x_{i-1})\}$.
- $F_i \cap F_j = \emptyset$, whenever $i \neq j$.

- Show that $\forall w \in \text{Pol}(A)$, w preserves R^P . Consider $\begin{pmatrix} f_1 \\ g_1 \end{pmatrix}, \dots, \begin{pmatrix} f_s \\ g_s \end{pmatrix} \in R^P$;

w applied to these tuples gives (f, g) where

$$f(x_1, \dots, x_p) = w(f_1(x_1, \dots, x_p), \dots, f_s(x_1, \dots, x_p))$$

$$g(x_1, \dots, x_p) = w(g_1(x_1, \dots, x_p), \dots, g_s(x_1, \dots, x_p))$$

Submaximal elements of \mathbb{P}

proof: • Choose one representative from every \sim -class \rightarrow denote this set of operations F_0 .

- Define $F_i = \{g \mid \exists f \in F_0 : g(x_1, \dots, x_p) = f(x_1, \dots, x_p, x_{i+1}, \dots, x_{i-1})\}$.
- $F_i \cap F_j = \emptyset$, whenever $i \neq j$.
- Show that $\forall w \in \text{Pol}(A)$, w preserves R^P . Consider $\begin{pmatrix} f_1 \\ g_1 \end{pmatrix}, \dots, \begin{pmatrix} f_s \\ g_s \end{pmatrix} \in R^P$;

w applied to these tuples gives (f, g) where

$$f(x_1, \dots, x_p) = w(f_1(x_1, \dots, x_p), \dots, f_s(x_1, \dots, x_p))$$

$$g(x_1, \dots, x_p) = w(g_1(x_1, \dots, x_p), \dots, g_s(x_1, \dots, x_p))$$

- $\Rightarrow f$ and g are from $\text{Pol}(A)$; moreover $g(x_1, \dots, x_p) = f(x_1, \dots, x_p, x_1)$
- $\Rightarrow w$ preserves $R^P \Rightarrow R^P$ is pp-definable over A .

Submaximal elements of \mathbb{P}

proof: • Choose one representative from every \sim -class \rightarrow denote this set of operations F_0 .

- Define $F_i = \{g \mid \exists f \in F_0 : g(x_1, \dots, x_p) = f(x_i, \dots, x_p, x_{i+1}, \dots, x_{i-1})\}$.
- $F_i \cap F_j = \emptyset$, whenever $i \neq j$.

- Show that $\forall w \in \text{Pol}(A)$, w preserves $R^{\mathbb{P}}$. Consider $\begin{pmatrix} f_1 \\ g_1 \end{pmatrix}, \dots, \begin{pmatrix} f_s \\ g_s \end{pmatrix} \in R^{\mathbb{P}}$;

w applied to these tuples gives (f, g) where

$$f(x_1, \dots, x_p) = w(f_1(x_1, \dots, x_p), \dots, f_s(x_1, \dots, x_p))$$

$$g(x_1, \dots, x_p) = w(g_1(x_1, \dots, x_p), \dots, g_s(x_1, \dots, x_p))$$

- $\Rightarrow f$ and g are from $\text{Pol}(A)$; moreover $g(x_1, \dots, x_p) = f(x_2, \dots, x_p, x_1)$
- $\Rightarrow w$ preserves $R^{\mathbb{P}}$ $\Rightarrow R^{\mathbb{P}}$ is pp-definable over A .

- $\mathbb{P} \leftrightarrow C_p : h: \mathbb{P} \rightarrow C_p \quad | \quad h': C_p \rightarrow \mathbb{P}; \text{choose } f_0 \in F_0; \text{all its cyclic shifts } f_i \in F_i$
 $F_i \mapsto i, f_i \quad | \quad i \mapsto f_i$

□

Submaximal elements of \mathbb{P}_3

Second step

Thm: Let A be a finite structure. A pp-contains $B_2 \Leftrightarrow \text{Pol}(A)$ does not satisfy Mal'cev.

proof: (similar) ...

Submaximal elements of \mathbb{P}_3

Second step

Thm: Let A be a finite structure. A pp-contains $B_2 \Leftrightarrow \text{Pol}(A)$ does not satisfy Mal'cev.

proof: (similar) ...

Next steps: Step 1 + Step 2 $\Rightarrow A$ has a 2-cyclic, 3-cyclic, and a Mal'cev operation.

BARTO, KOZIK
 \Rightarrow p -cyclic, if p prime.

Submaximal elements of \mathbb{P}_3

Second step

Thm: Let A be a finite structure. A pp-contains $B_2 \Leftrightarrow \text{Pol}(A)$ does not satisfy Mal'cev.

proof: (similar) ...

Next steps: \triangleright Step 1 + Step 2 $\Rightarrow A$ has a 2-cyclic, 3-cyclic, and a Mal'cev operation.

BARTO, KOZIK
 \Rightarrow p -cyclic, $\nparallel p$ prime.

\triangleright p -cyclic, $\nparallel p$ + Mal'cev $\Rightarrow A$ has a majority operation \rightsquigarrow symmetric maj.

Submaximal elements of \mathbb{P}_3

Second step

Thm: Let A be a finite structure. A pp-contains $B_2 \Leftrightarrow \text{Pol}(A)$ does not satisfy Mal'cev.

proof: (similar) ...

Next steps: ▶ Step 1 + Step 2 $\Rightarrow A$ has a 2-cyclic, 3-cyclic, and a Mal'cev operation.

BARTO, KOZIK
▶ p -cyclic, $\nparallel p$ prime.

- ▶ p -cyclic, $\nparallel p$ + Mal'cev $\Rightarrow A$ has a majority operation \rightsquigarrow symmetric maj.
- ▶ majority + Mal'cev $\Rightarrow A$ has a minority operation \rightsquigarrow symmetric minority.

Submaximal elements of \mathbb{P}_3

Second step

Thm: Let A be a finite structure. A pp-contains $B_2 \Leftrightarrow \text{Pol}(A)$ does not satisfy Mal'cev.

proof: (similar) ...

Next steps: ▶ Step 1 + Step 2 $\Rightarrow A$ has a 2-cyclic, 3-cyclic, and a Mal'cev operation.

BARTO, KOZIK
▶ p -cyclic, $\nvdash p$ prime.

- ▶ p -cyclic, $\nvdash p$ + Mal'cev $\Rightarrow A$ has a majority operation \rightsquigarrow symmetric maj.
- ▶ majority + Mal'cev $\Rightarrow A$ has a minority operation \rightsquigarrow symmetric minority.
- ▶ symm. maj + minor. + 2-cyclic $\Rightarrow A$ has a totally symm. operation of arity n , $\forall n \geq 2$.

Submaximal elements of P_3

Second step

Thm: Let A be a finite structure. A pp-contains $B_2 \Leftrightarrow \text{Pol}(A)$ does not satisfy Mal'cev.

proof: (similar) ...

Next steps: Step 1 + Step 2 $\Rightarrow A$ has a 2-cyclic, 3-cyclic, and a Mal'cev operation.

BARTO, KOZIK
 $\Rightarrow p\text{-cyclic, } \forall p \text{ prime.}$

- $p\text{-cyclic, } \forall p \oplus \text{Mal'cev} \Rightarrow A$ has a majority operation \rightsquigarrow symmetric maj.
- majority \oplus Mal'cev $\Rightarrow A$ has a minority operation \rightsquigarrow symmetric minority.
- symm. maj \oplus minor. \oplus 2-cyclic $\Rightarrow A$ has a totally symm. operation of arity $n, \forall n \geq 2$.

$$S(x_1, \dots, x_n) := \text{tup} \left(S_{n-1}(x_1, M^c(x_1, x_2, x_3), x_4, \dots, x_n), \right. \\ \left. S_{n-1}(x_2, M^c(x_1, x_2, x_3), x_4, \dots, x_n), \right. \\ \left. S_{n-1}(x_1, M^c(x_1, x_2, x_3), x_4, \dots, x_n) \right)$$

Submaximal elements of P_3

The addition

Def: An **addition** of arity n is an idempotent operation $m_n^c : A^n \rightarrow A$ s.t.

$$m_n^c(x_1, \dots, x_n) = \begin{cases} c \in A & \\ \text{the only value occurring} \\ \text{an odd number of times} & \end{cases}$$

if there are at least three different values occurring in the tuple (x_1, \dots, x_n)
otw.

Submaximal elements of P_3

The addition

Def: An **addition** of arity n is an idempotent operation $m_n^c : A^n \rightarrow A$ s.t.

$$m_n^c(x_1, \dots, x_n) = \begin{cases} c \in A & \\ \text{the only value occurring} \\ \text{an odd number of times} & \end{cases}$$

if there are at least three different values occurring in the tuple (x_1, \dots, x_n)
otw.

► D. Zhu $\Rightarrow m_{2n+1}^c := m_3^c (\overbrace{\quad}^1, \overbrace{\quad}^1)$

Submaximal elements of P_3

The addition

Def: An **addition** of arity n is an idempotent operation $m_n^c : A^n \rightarrow A$ s.t.

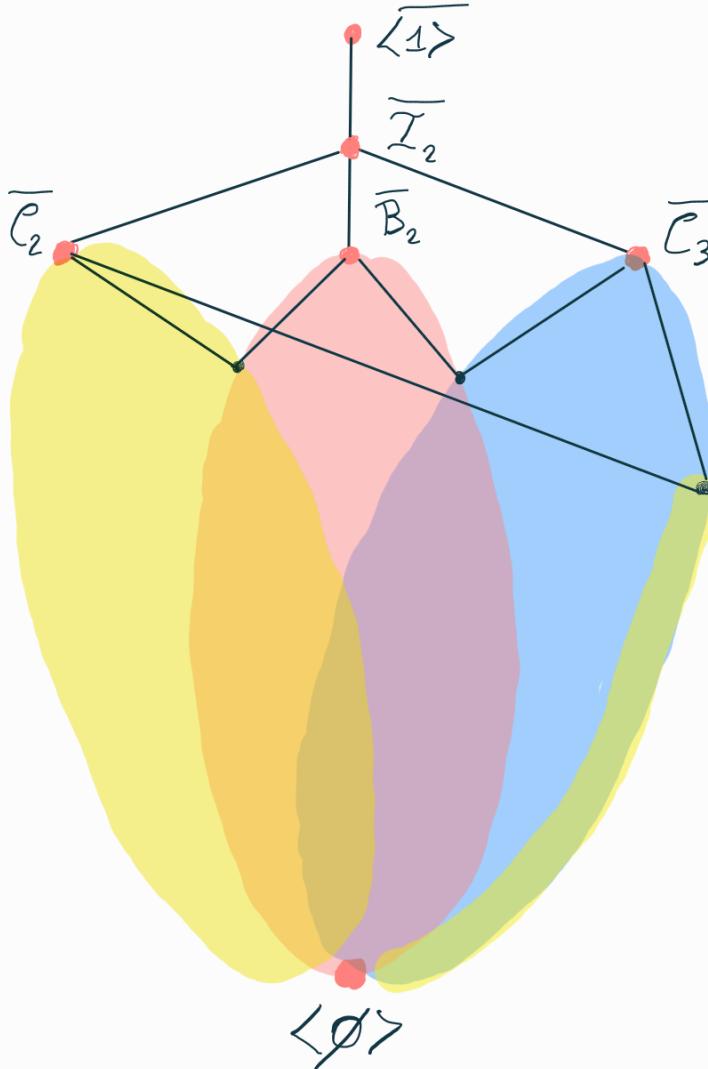
$$m_n^c(x_1, \dots, x_n) = \begin{cases} c \in A & \text{if there are at least three different values occurring in the tuple } (x_1, \dots, x_n) \\ \text{the only value occurring an odd number of times} & \text{otw.} \end{cases}$$

► D. Zhu $\Rightarrow m_{2n+1}^c := m_3^c (\overbrace{\quad}, \overbrace{\quad}, \overbrace{\quad})$

Thm: There is a minor preserving map from \mathbb{F}_2 to A (where A satisfies \oplus).

\sum : sum of l monomials $\mapsto m_l(s_{n_1}(\dots), \dots, s_{n_l}(\dots))$

P₃ : an overview



- : fully described [B V Z]
- : potentially 2^ω elements
- : work in progress
[FORAVANTI, ROSSI, V.]

A hand-drawn style "Thank You!" message. The word "Thank" is written in dark blue cursive, and "You" is written in black cursive. A large, multi-layered oval shape overlaps the letters, containing segments of yellow, orange, pink, purple, and blue. Above the ovals, a small black line drawing of a tree with red dots at its branches and trunk is positioned above the letter "T". To the right of the "You" is a large, thick black exclamation mark.

Thank You !